

Fiddly Lemma

\mathcal{C}, \mathcal{D} ab cats, $F: \mathcal{C} \rightarrow \mathcal{D}$ additive

For any $\varphi: X \rightarrow Y$ in \mathcal{C} , \exists map
 $F(\ker \varphi) \rightarrow \ker(F\varphi)$.

Lemma F is left-exact

\Leftrightarrow this map is an iso $\forall \varphi$.

Proof Assume F left exact

For any $\varphi: X \rightarrow Y$

$$0 \rightarrow \ker \varphi \rightarrow X \rightarrow \operatorname{Im} \varphi \rightarrow 0$$

$$0 \rightarrow \operatorname{Im} \varphi \rightarrow Y \rightarrow \operatorname{coker} \varphi \rightarrow 0$$

both exact.

Hence $0 \rightarrow F(\ker \varphi) \rightarrow FX \rightarrow F(\operatorname{Im} \varphi)$
& $0 \rightarrow F(\operatorname{Im} \varphi) \rightarrow FY \rightarrow F(\operatorname{coker} \varphi)$
exact.

2nd seq says that $F(\operatorname{Im} \varphi) \rightarrow FY$ is
a monomorphism

Hence $\ker(FX \rightarrow F(\operatorname{Im} \varphi))$

and $\ker(FX \rightarrow FY)$ are the same.

so exactness of $F(\text{sequence 1})$
gives

$$0 \rightarrow F(\ker \varphi) \rightarrow FX \xrightarrow{F\varphi} FY$$

exact, ie $\ker(F\varphi) = F(\ker \varphi)$.

Converse implication: exercise. \square

Apply this to $F^{op}: \mathcal{C}^{op} \rightarrow \mathcal{D}^{op}$

to see that right-exact \Leftrightarrow
preserves cokernels.

Hence exact functors preserve all
exact seqs (not just short ones).

Notation If $F: \mathcal{C} \rightarrow \mathcal{D}$ contravariant,

say F is left exact

if $F: \mathcal{C}^{op} \rightarrow \mathcal{D}$ left exact.

Prop If \mathcal{C} ab. cat, $X \in \text{Ob}(\mathcal{C})$

the functors $\text{Hom}(X, -): \mathcal{C} \rightarrow \underline{\text{Ab}}$

$\text{Hom}(-, X): \mathcal{C} \rightarrow \underline{\text{Ab}}$ (contra)

are left exact.

Proof suffices to prove 1st claim

Let $0 \rightarrow Y \xrightarrow{\alpha} Y' \xrightarrow{\beta} Y''$ exact in \mathcal{C}

Need to show: i) if $f \in \text{Hom}(X, Y)$ st $\alpha \cdot f = 0$

then $f = 0$. This is exactly
claim that α is a monomorphism. ✓

ii) if $f: X \rightarrow Y'$ st $\beta \cdot f = 0$ then

$\exists!$ lifting to $X \rightarrow Y$.

= universal property of $\text{Ker}(\beta) = Y$. ✓

□

This gives lots of left-exact functors.

Key Example G group,

\mathcal{C} = cat of G -modules

[abelian gps with left G -action
that is \mathbb{Z} -linear]

Claim $\mathcal{C} \rightarrow \text{Ab}$

$M \mapsto M^G$ invariants functor

is left-exact.

PF consider $\mathbb{Z} \in \text{Obj}(\mathcal{C})$ with G acting

trivially
Have $\text{Hom}_{\mathcal{C}}(\mathbb{Z}, -) = (-)^G$

so $(-)^G$ is left exact.

It is not exact: eg let $G =$ infinite cyclic

$0 \rightarrow \left(\begin{array}{c} \mathbb{Z} \text{ with} \\ G \text{ triv} \end{array} \right) \rightarrow \left(\begin{array}{c} \mathbb{Z}^2 \text{ with} \\ g \text{ acting} \\ \text{as } \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \end{array} \right) \xrightarrow{\langle g \rangle} \mathbb{Z} \rightarrow 0$

short exact in \mathcal{C} .

$0 \rightarrow \mathbb{Z} \xrightarrow{\text{id}} \mathbb{Z} \xrightarrow{\circ} \mathbb{Z}$

not exact at RH end.

(will understand this later
using group cohomology.)

Lemma let \mathcal{C}, \mathcal{D} ab. cats

$L: \mathcal{C} \rightarrow \mathcal{D}$ additive, $R: \mathcal{D} \rightarrow \mathcal{C}$
right adjt to L

Then L is right exact, & R is left exact.

PF R preserves all limits (see Sheet 1)

\Rightarrow preserves kernels.

similarly L preserves all colimits.

Corollary $R \xrightarrow{\varphi} S$ morphism of rings

R -Mod \rightarrow S -Mod

$M \mapsto S \otimes_{R, \varphi} M.$

This functor is right exact, being
left adj^t to the functor

$\varphi^*: \underline{S\text{-Mod}} \rightarrow \underline{R\text{-Mod}}.$

§ 2.4 Projective & injective objs

\mathcal{C} ab. cat.

Defⁿ Say $X \in \text{Obj}(\mathcal{C})$ is
projective if $\text{Hom}(X, -)$ exact
injective if $\text{Hom}(-, X)$ " .

Equivalently:

X is proj. if "morphisms from X to
quotients always lift"

i.e. if $Y \twoheadrightarrow Y'$, any $X \rightarrow Y'$
lifts to Y .

X is inj if "morphisms from subobjects
to X are extendable."

Basic example: will see that

in Ab, \mathbb{Z} is proj,

\mathbb{Q}/\mathbb{Z} is inj.

Defⁿ \mathcal{C} has enough projectives

if every $X \in \text{Ob}(\mathcal{C})$ has a surjⁿ

$P \twoheadrightarrow X$ with P proj.

enough injectives if every X has

$X \hookrightarrow I$, I inj.

R ring.

Thm (i) In R -Mod, free modules

$$R^{(\Sigma)} = \bigoplus_{\sigma \in \Sigma} R \quad (\Sigma \text{ any set})$$

are proj.

(ii) R -Mod has enough proj.

(iii) in R -Mod, M is proj $\iff \exists$ module N

st $M \oplus N$ is free.

Proof i) Suppose $F = R^{(\Sigma)}$ is free.

$Y \twoheadrightarrow Y'$ epi. ^(*) For each $\sigma \in \Sigma$,

lift ϕ 's th gen. of F to Y arbitrarily.

$(*) \phi: F \rightarrow Y'$ This gives a lifting

$$\hat{\phi}: F \rightarrow Y.$$

(ii) For any M , free module on underlying set of M surjects onto M .

(NB: This also shows that $(R\text{-Mod})^{\text{fg}}$,
cat. of finitely gen. R -mods, has enough
projs; it does not have enough inj.s.)

(iii) Suppose P proj and take a surjⁿ
 $F \twoheadrightarrow P$, F free. Then

$\text{id}: P \rightarrow P$ lifts to a map $P \xrightarrow{\alpha} F$,
and can check that $\text{im}(\alpha) \cong P$
and $F = \ker(\beta) \oplus \text{im}(\alpha)$.

Conversely if $F = X \oplus X'$ with F free,
 $Y \twoheadrightarrow Y'$, $\varphi: X \rightarrow Y'$. Then composite

$F \twoheadrightarrow X \xrightarrow{\varphi} Y'$ lifts to $F \rightarrow Y$
& compose it with $X \hookrightarrow F$
to lift φ .

Much harder theorem:

Thm $R\text{-Mod}$ has enough injectives.

Ideas of pf: ① reduce to $R = \mathbb{Z}$
② use some kind of "duality"

Define $\Omega =$ the R -module $\text{Hom}_{\mathbb{Z}}(R, \mathbb{Q}/\mathbb{Z})$
with $(r \cdot \varphi)(-) = \varphi((-) \cdot r)$.

Lemma 1 For any R -module M , $\text{Hom}_R(M, \Omega)$
 $\cong \text{Hom}_{\mathbb{Z}}(M, \mathbb{Q}/\mathbb{Z})$ (naturally in M).

Proof Given $\varphi: M \rightarrow \mathbb{Q}/\mathbb{Z}$ group hom,

let $\tilde{\varphi}: M \rightarrow \Omega$ be def by
 $m \mapsto (r \mapsto \varphi(rm))$

conversely φ is image of $\tilde{\varphi}$ under evalⁿ
check these are inverse operations. at $1 \in R$.

Lemma 2 Ω is injective.

Pf By Lemma 1 suffices to show
that \mathbb{Q}/\mathbb{Z} is inj. in Ab.

Induction step: if $G \supset H$ ab gps,
 $\varphi: H \rightarrow \mathbb{Q}/\mathbb{Z}$, G/H cyclic gen. by some g .

If g has ∞ order in G/H , then $G \cong H \oplus \mathbb{Z}$
and can extend φ to G by letting $\varphi(g)$ be
arbitrary.

If g has order $n < \infty$, then $\varphi(g^n) \in \mathbb{Q}/\mathbb{Z}$
is given, & since \mathbb{Q}/\mathbb{Z} is divisible can
define a hom $G \rightarrow \mathbb{Q}/\mathbb{Z}$ by choosing
 $\varphi(g)$ st $n \cdot \varphi(g) = \varphi(g^n)$.

Zorn's lemma \Rightarrow our lemma follows.

Lemma 3 For any R -module M , and
 $m \neq 0 \in M$, \exists hom $M \xrightarrow{\varphi} \mathbb{Q}$
st $\varphi(m) \neq 0$.

Pf Again can assume $R = \mathbb{Z}$ wlog by Lem 1.

Since \mathbb{Q}/\mathbb{Z} inj, can assume M generated
by m . If m has order n send it to $\frac{1}{n}$
 ∞ , (anything $\neq 0$).

P.

Pf of Thm

Define $D(M)$, for M in $R\text{-Mod}$,
as $\text{Hom}_R(M, \Omega)$ (with some $R\text{-mod}$
str.)

For any M have nat^l map

$$M \rightarrow D(D(M))$$
$$m \mapsto (\varphi \mapsto \varphi(m))$$

By L3 this is an injection.

Will show $D(D(M))$ injects into an injective.

Take surjⁿ $F \rightarrow D(M)$ $F \cong R^{(\Sigma)}$ free

thus get map $D(D(M)) \rightarrow D(F)$

and $D(F)$ is a direct sum of
copies of Ω , hence injective. \square

3. Complexes and Homotopies

3.1 The category of complexes

\mathcal{C} ab. cat.

Defⁿ $\text{Ch}(\mathcal{C}) =$ category whose objs are cochain complexes in \mathcal{C} , $(X^i)_{i \in \mathbb{Z}}$
+ morphisms being collections $(f^i: A^i \rightarrow B^i)_{i \in \mathbb{Z}}$
compat. w. the differentials

$$\begin{array}{ccc} A^i & \xrightarrow{d_A^i} & A^{i+1} \\ f^i \downarrow & & \downarrow f^{i+1} \\ B^i & \xrightarrow{d_B^i} & B^{i+1} \end{array} \quad \text{commutes } \forall i.$$

Prop (easy): $\text{Ch}(\mathcal{C})$ is an ab. cat. with obvious notions of kernel, cokernel.

Defⁿ the "nth cohomology functor"

$\text{Ch}(\mathcal{C}) \rightarrow \mathcal{C}$ sends A^i to $H^n(A^i)$

Problem: not actually well-def! - kernels + cokernels not unique. Want to pick a ker & coker for every morphism in \mathcal{C} .

In practice most ab cats come with "standard choices" of ker/coker.

Defⁿ A morphism f in $\mathcal{C}h(\mathcal{C})$ is a quasi-isomorphism if it induces isomorphisms on H^i for all $i \in \mathbb{Z}$.

Lemma ("Snake Lemma") for any diagram

$$\begin{array}{ccccccc}
 \text{Ker } f & \xrightarrow{\quad} & \text{Ker } g & \xrightarrow{\quad} & \text{Ker } h & & \\
 \downarrow & & \downarrow & & \downarrow & & \\
 A & \longrightarrow & B & \longrightarrow & C & \longrightarrow & 0 \\
 \downarrow f & & \downarrow g & & \downarrow h & & \\
 0 \longrightarrow & A' & \longrightarrow & B' & \longrightarrow & C' & \\
 \downarrow \text{coker } f & & \downarrow \text{coker } g & & \downarrow \text{coker } h & & \\
 & \longrightarrow & & \longrightarrow & & &
 \end{array}$$

in \mathcal{C} , commutative with exact rows,

\exists arrow $\text{Ker}(h) \rightarrow \text{Coker}(f)$

st $\text{ker}(f) \rightarrow \text{ker}(g) \rightarrow \text{ker}(h) \rightarrow \text{coker } f \rightarrow \dots$
is exact.

Can prove this directly for $R\text{-Mod}$, see next time.

For gen^l \mathcal{C} , can reduce to this via Freyd-Mitchell embedding.