

Outer space and Automorphisms of free groups

LECTURE 7 Change of time next week

Last time, showed how to compute
 $H_*(\text{Out}F_n; \mathbb{Q}) = H_*(K_n/\text{Out}F_n; \mathbb{Q})$

using a chain complex derived from the
cube complex structure of K_n

Today: connection of $H^*(K_n/\text{Out}F_n; \mathbb{Q})$ with
Kantsevich's graph homology.

All homology will be with trivial coefficients
in a field of char. 0, so we will omit the
coefficients in the notation

$H_*(K_n / \text{cut } F_n)$ is the homology of the chain complex C_* with $C_k = \bigoplus_{(G, \Phi)} \mathbb{Q}$

where • G connected, $\pi_1 G \cong F_n$, all vertices have valence ≥ 3

• Φ is a forest in G with k_2 edges e_1, \dots, e_k

• (G, Φ) has no orientation-reversing automorphisms

and • $\partial(G, \Phi) = \sum_{i=1}^k (-1)^{i+1} (G, \Phi - e_i) + (-1)^i (G_{e_i}, \Phi_{e_i})$

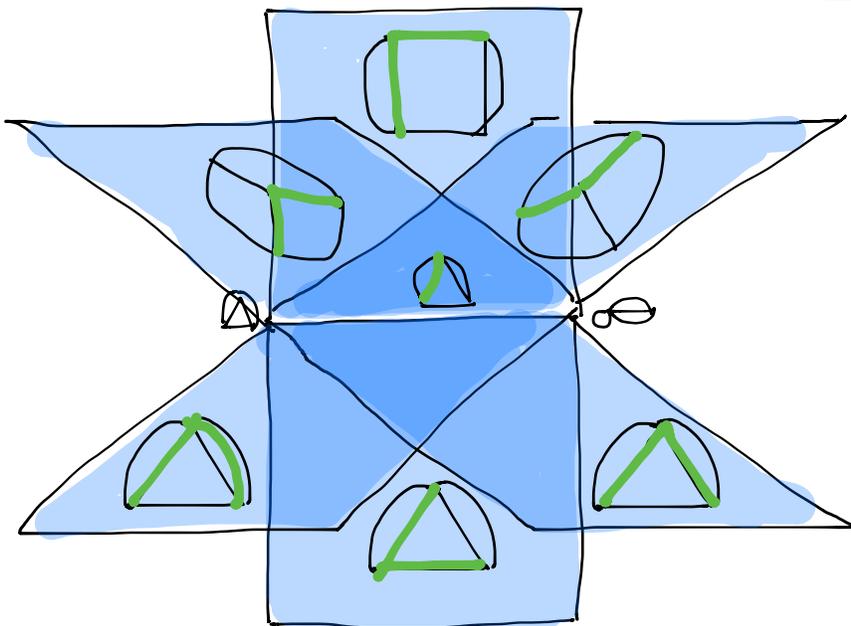
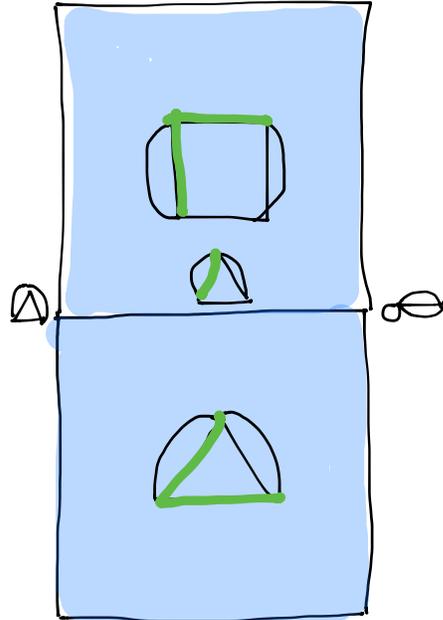
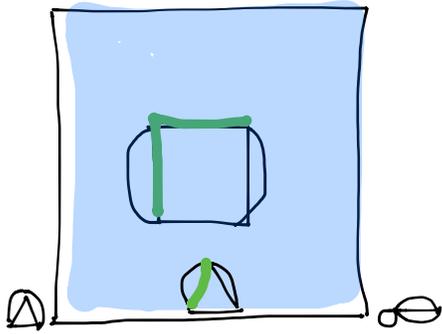
$$= \partial_R(G, \Phi) + \partial_C(G, \Phi)$$

∂_R "removes edges from Φ "

∂_C "collapses edges of Φ "

To make the connection with Kontsevich's graph homology, need to simplify our chain complex. And, turns out it's more convenient to use the cochain complex.

Coboundary of $\sigma = \triangle \xrightarrow{\quad} \circlearrowleft$



Cochain complex A cube (G, Φ) is

a codimension 1 face of cubes

$$(G, \Phi \cup e) \text{ and } (G^\alpha, \Phi^\alpha)$$

where. e is any edge which can be added to Φ

without forming a cycle

α blows up some vertex of G into 2 vertices, and Φ^e contains the new edge.

$$\begin{aligned} \delta(G, \Phi) &= \delta_R(G, \Phi) + \delta_C(G, \Phi) \\ &= \sum_e (G, \Phi \cup e) + \sum_\alpha (G^\alpha, \Phi^\alpha) \end{aligned}$$

where $\Phi \cup e$ and Φ^α are oriented appropriately

In sphere system model, $G \rightarrow \Delta_G \subset M_n$
 Δ_G a simple sphere system, spheres \leftrightarrow edges of G
 Φ a forest \leftrightarrow sub-system $\Delta_\Phi \subset \Delta_G$ (not simple)
 $(G, \Phi) \leftrightarrow (\Delta_G, \Delta_\Phi)$

Then

$$(G^\alpha, \Phi^\alpha) \leftrightarrow (\Delta_G \cup \Delta_\alpha, \Delta_\Phi \cup \Delta_\alpha)$$

$$\text{so } \delta_c(G, \Phi) = \delta_c(\Delta_G, \Delta_\Phi)$$

$$= \sum_{\substack{\Delta_\alpha \text{ compatible} \\ \text{with } \Delta_G}} (\Delta_G \cup \Delta_\alpha, \Delta_\Phi \cup \Delta_\alpha)$$

We have

$$C_k = \bigoplus_{\substack{(G, \Phi) \\ e(\Phi) = k}} \mathbb{Q} = \bigoplus_{e(\Phi) = k} (G, \Phi) \quad [\text{Notation}]$$

We can decompose this further by # vertices of G

$$C_k = \bigoplus_{v > k} C_{v, k}$$

$$\text{where } C_{v, k} = \bigoplus_{\substack{e\Phi = k \\ vG = v}} (G, \Phi)$$

Top-dimil cubes are

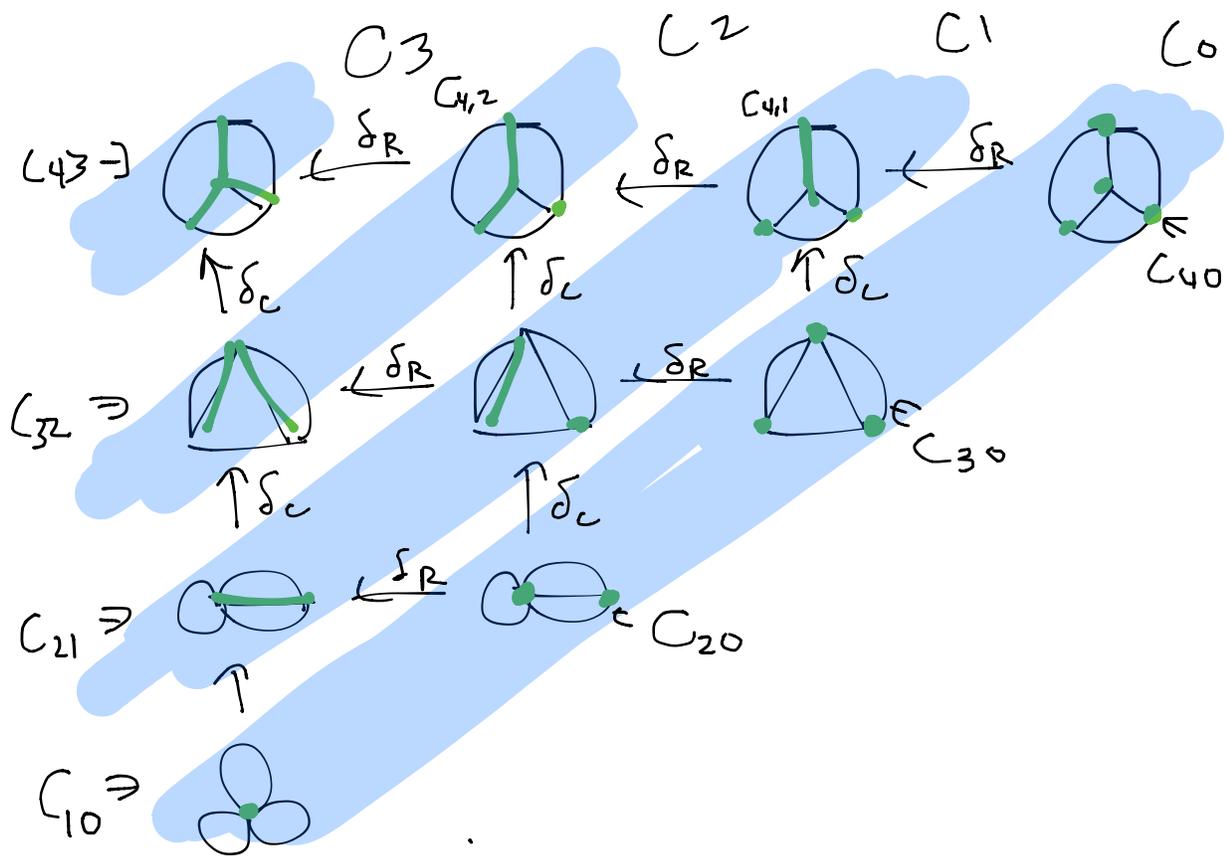
$$(G, T) = (\text{trivalent graph, maximal tree}) \\ (vG = 2n - 2, eT = 2n - 3)$$

$$\text{next are } v(G) = 2n - 2, e\bar{\Phi} = 2n - 4$$

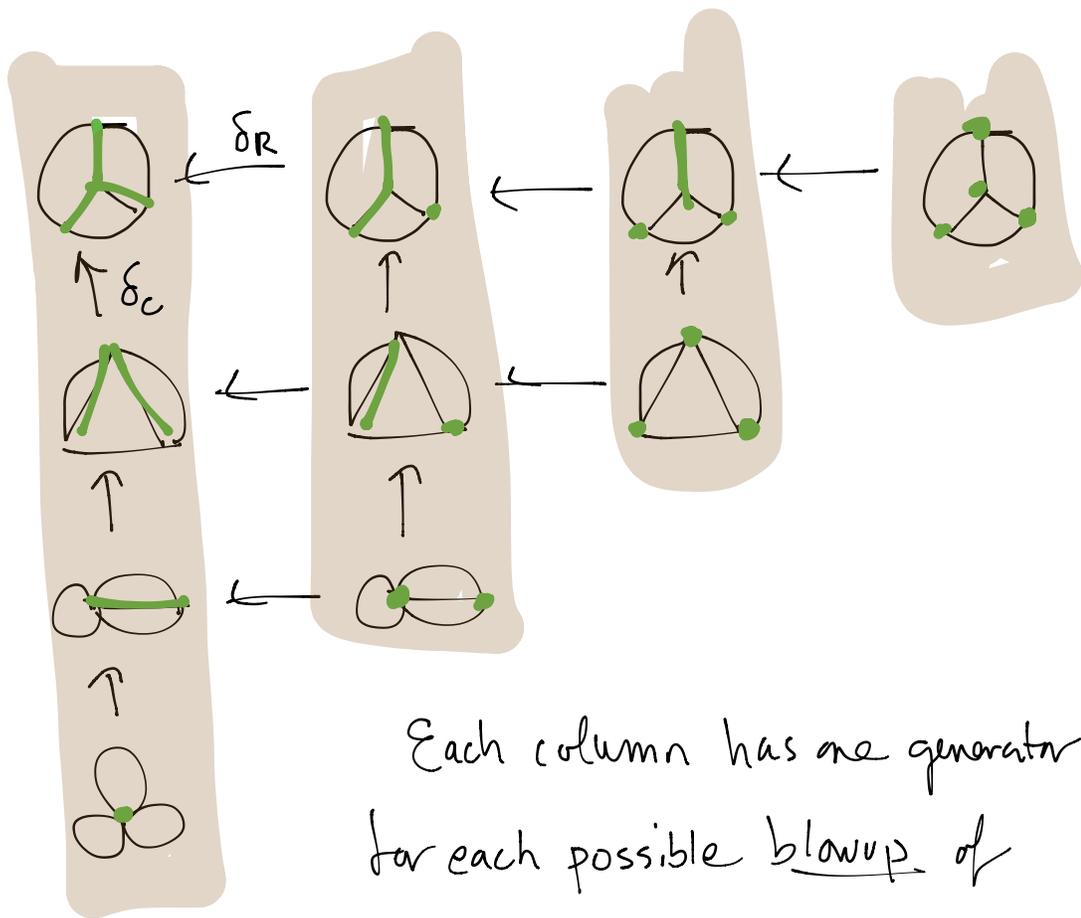
$$\text{and } vG = 2n - 3, e\bar{\Phi} = 2n - 4$$

Arrange the $C_{v, k}$ in a grid

Draw a **sample** (G, Φ) at each spot to keep this straight:



We can compute $H^*(C_*)$ by first computing the vertical homology, then the horizontal H_* (this is spectral sequences, but don't tell anybody)



Each column has one generator for each possible blowup of the graphs G at the bottom

Different blowups may result in isomorphic graphs, e.g. if G has an automorphism

If these were marked graphs, then the complex of blowups is $\mathbb{A}e_{>G}$ in K_n , which we know is $\cong VS^{2n-3-r(G)}$

Since they are unmarked graphs, we need to quotient \mathbb{R}^k/G by its stabilizer, which is $\cong \text{Aut}(G)$

\mathbb{R}^k/G has no homology except in the top dimension

So same is true for $\mathbb{R}^k/G/\text{Aut}(G) \subset \mathbb{R}^n/\text{Aut}(F_n)$

[pf uses equivariant H_* spectral sequence:

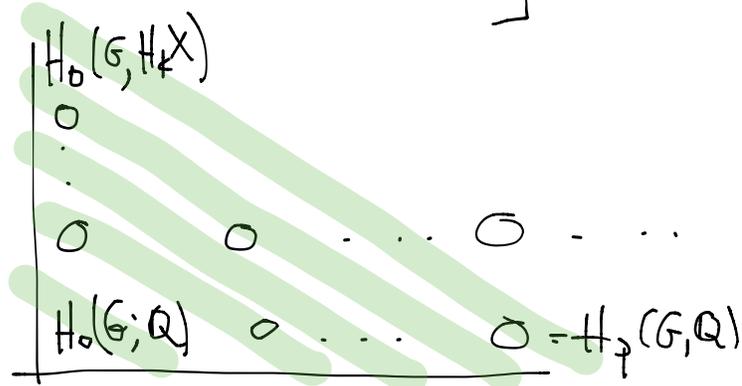
$X \cong VS^k$, G finite acting on X cellularly

$$E_{p,q}^2 = H_p(G; H_q(X)) \Rightarrow H_{p+q}(X/G)$$

Since $H_q(X) = 0$ unless $q = 0$ or k and

$H_p(G; \mathbb{Q}) = 0$ for $p > 0$, get $H_{p+q}(X/G) = 0$ for $0 < p+q < k$

(also $H_{p+q}(X/G) = 0$ for $p+q > k$)



So after taking vertical H^* , co-chain complex becomes

$$0 \leftarrow \bigoplus_{e(T)=2n-2} (G, T) / \text{Im } \delta_c \xleftarrow{\delta_R} \bigoplus_{e(\Phi)=2n-3} (G, \Phi) / \text{Im } \delta_c \xleftarrow{\delta_R} \dots$$

Where G is always trivalent.

$$\leftarrow \dots \bigoplus_{e(\Phi)=0} (G, -)$$

Much smaller chain complex!

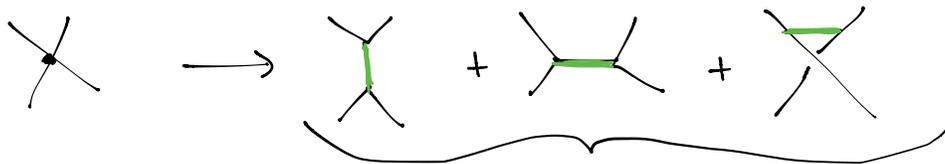
$$\delta_c : C_{2n-3, k-1} \rightarrow C_{2n-2, k}$$

sends $(G, \Phi) \rightarrow \sum_{\alpha} (G^{\alpha}, \Phi^{\alpha})$

G^{α} trivalent \Rightarrow

G is trivalent except at one 4-valent vertex v 

There are 3 ways of expanding this into two trivalent vertices. The new edge becomes part of the forest.



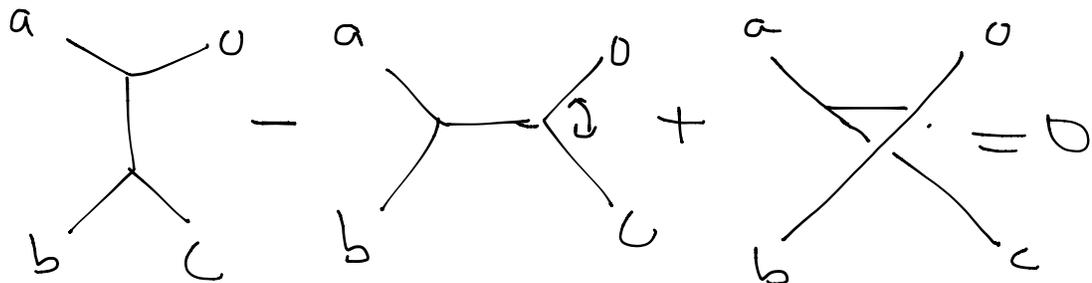
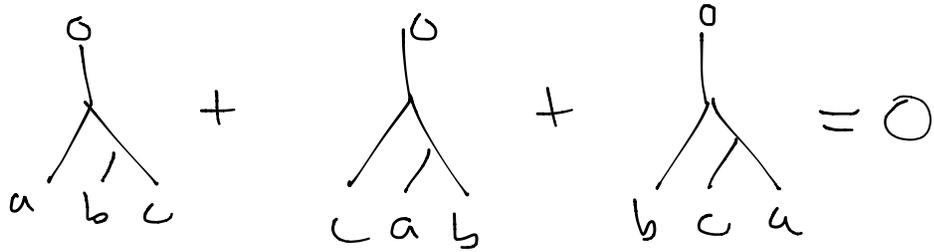
This is called an IHX-vector.

So our chain complex is

$$0 \leftarrow \frac{\bigoplus_{\substack{e\Phi=2n-3}} (G, \mathcal{T})}{\text{IHX}} \leftarrow \frac{\bigoplus_{e\Phi=2n-4} (G, \mathcal{T}_e)}{\text{IHX}} \leftarrow \dots$$

This is the first hint that $H^*(\text{Out}F_n)$ is related to Lie algebras: (actually the Lie operad)
 Modding out by IHX is an encoding of the Jacobi identity:

$$[a, [b, c]] + [c, [a, b]] + [b, [c, a]] = 0$$



Kontsevich's theory

We now turn from something easy to understand (graphs) to something harder to understand (algebra!). Aim: convert algebras into graphs.

$V_d = \mathbb{R}^{2d}$ equipped with standard symplectic

form $\langle u, w \rangle$, ie

\mathbb{R}^{2d} has basis $B = \{p_1, \dots, p_d, q_1, \dots, q_d\}$

$\langle p_i, q_i \rangle = 1 = -\langle q_i, p_i \rangle$ all other $\langle x, y \rangle = 0$

for $x, y \in B$.

$C_d =$ free commutative algebra on B
= polynomials in the p_i, q_i with real coeffs.

$A_d =$ free associative algebra on B
= polynomials in non-commuting variables

p_i, q_i

$L_d =$ free Lie algebra on B : generated by
bracket expressions in the p_i, q_i

(eg $[p_1, p_2]$ or $[[p_1, [q_2, p_1]], [q_2, q_4]]$)

modulo antisymmetry $[X, Y] = -[Y, X]$

Jacobi identity

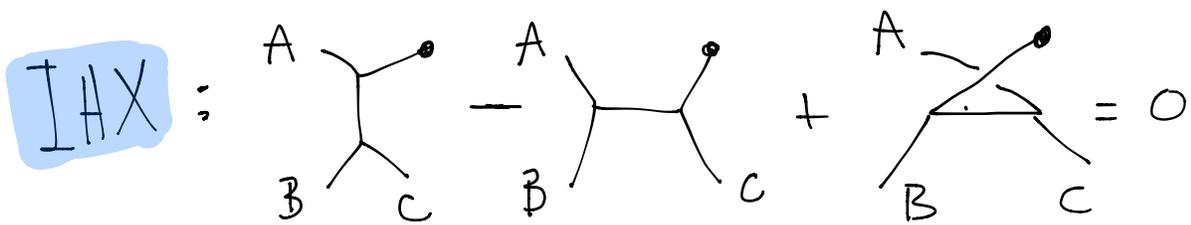
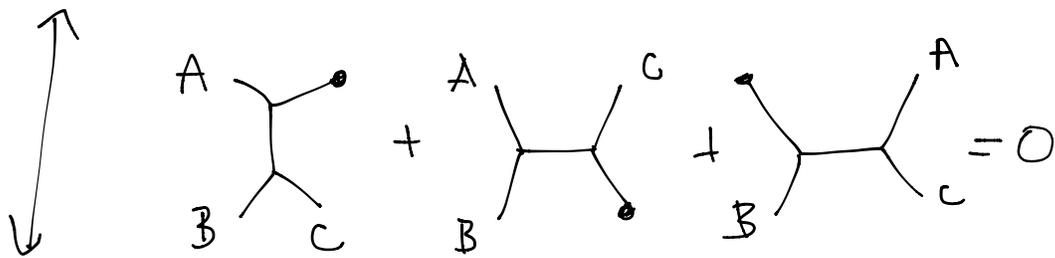
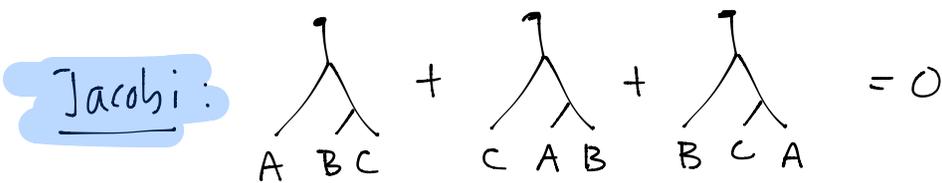
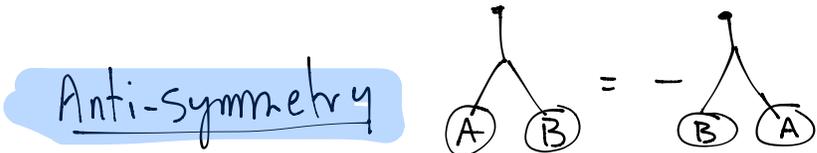
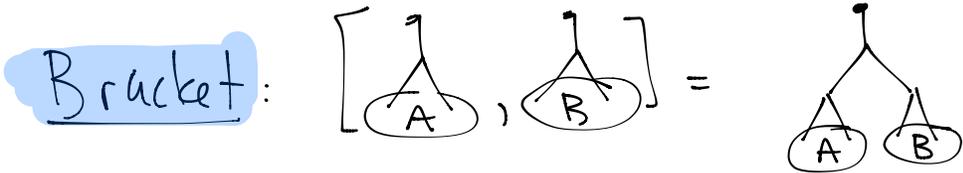
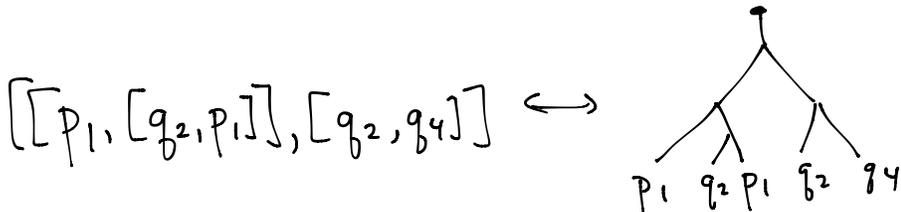
$$[A, [B, C]] + [C, [A, B]] + [B, [C, A]] = 0$$

Now turn this into algebras of labelled, rooted

graphs: Start in L_d

Ld:

bracket expression \leftrightarrow rooted planar binary tree



Multi-Linearity

$$\begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ c+x \quad y \quad z \quad w \end{array} = c \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ x \quad z \quad w \end{array} + \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ y \quad z \quad w \end{array}$$

A_d generated by planar rooted labelled star-graphs:

$$\begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ p_1 \quad q_2 \quad p_1 \quad q_3 \end{array} \leftrightarrow \text{monomial } p_1 q_2 p_1 q_3$$

Product: $\begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ a \quad b \quad c \end{array} \cdot \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ d \quad e \end{array} = \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ a \quad b \quad c \quad d \quad e \end{array}$

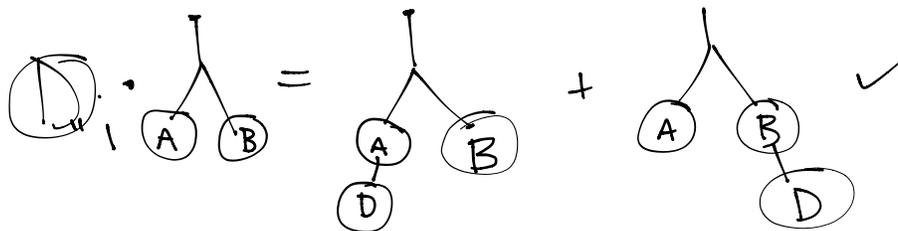
C_d : generated by rooted trees (no planar embedding)

(so $\begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ x \quad y \quad z \end{array} = \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ x \quad z \quad y \end{array} = \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ y \quad x \quad z \end{array} = \dots$)

Product: $\begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ x \quad y \quad z \end{array} \cdot \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ u \quad v \end{array} = \begin{array}{c} \bullet \\ | \\ \diagup \quad | \quad \diagdown \\ x \quad y \quad z \quad u \quad v \end{array}$

ie: sum $\langle x, a \rangle$ = Result of "hanging D off T"
 over all ways of pairing a leg x
 of D with a leg a of T

It's clear that D is a derivation:



Exercise:
$$\begin{matrix} p_1 \\ | \\ p_2 \quad q_1 \end{matrix} \cdot \left(\sum_{i=1}^d \begin{matrix} \cdot \\ | \\ p_i \quad q_i \end{matrix} \right) = 0$$

More ambitious exercise:
$$w = \sum_{i=1}^d \begin{matrix} \cdot \\ | \\ p_i \quad q_i \end{matrix}$$

then $Dw = 0$ for any $D \in \mathfrak{h}_d$.

Exercise: What does "hanging D onto T"
 mean for $H_d = A_d$ or C_d ?

Another definition of \mathfrak{h}_d
 = derivations of H_d killing ω ,
 where $\omega = \sum [p_i, q_i]$ if $H_d = L_d$
 $\omega = \sum p_i q_i - q_i p_i$ if $H_d = A_d$
 $\omega = \sum dp_i \wedge dq_i$ if $H_d = C_d$

Derivations of an algebra form a Lie algebra
 by defining $[D, D'] = D \circ D' - D' \circ D$

We can describe the bracket using graphs
 as follows:

Let $\Delta_i = \begin{array}{c} x \\ \diagup \quad \diagdown \\ \text{---} \text{---} \text{---} \\ \diagdown \quad \diagup \\ u \quad z \end{array}$, $\Delta'_i = \begin{array}{c} x' \\ \diagup \quad \diagdown \\ \text{---} \text{---} \text{---} \\ \diagdown \quad \diagup \\ z' \end{array}$ be spiders
 in H_d . Given a leg $x \in \Delta_i$ and $x' \in \Delta'_i$,
 can mate Δ_i and Δ'_i by gluing x_i to x'_i ,
 contracting the new edge if $H = A$ or C , and

multiplying by $\langle x_i, x_i' \rangle$

Then $[s_i, s_i']$ is the sum of all possible matings.

eg in L_d :

$$\left[\begin{array}{c} q_1 \\ \diagdown \quad \diagup \\ p_1 \quad p_2 \end{array}, \begin{array}{c} q_1 \\ \diagdown \quad \diagup \\ p_3 \quad p_1 \end{array} \right] = - \begin{array}{c} q_1 \\ \diagdown \quad \diagup \\ p_3 \quad p_1 \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad p_2 \end{array} + \begin{array}{c} q_1 \\ \diagdown \quad \diagup \\ p_3 \quad p_1 \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad p_1 \quad p_2 \end{array}$$

(all other terms have coefficient 0)

$[A, B] = -[B, A]$ since each coefficient

$\langle x, x' \rangle$ becomes $\langle x', x \rangle = -\langle x, x' \rangle$

Jacobi identity is also true, but requires more effort to verify

The natural embedding $V_d \hookrightarrow V_{d+1}$ induces inclusions $h_d \hookrightarrow h_{d+1}$ (we're just allowing more labels). Define $h_\infty = \varinjlim h_d$

Kontsevich's Theorem:

$H_d = L_d$. Then

" $H_* l_\infty$ computes $H^*(\text{Out } F_n)$ for all n "

More precisely:

$$P H_* l_\infty \cong H^*(sp_\infty) \oplus \bigoplus_{n \geq 2} H^*(\text{Out } F_n)$$

$H_* l_\infty$ is a Hopf algebra,
this is the "primitive part"

l_∞ contains a copy
of sp_∞ , corresp.
to 2-legged
spiders