

# Lecture 5 Graph homology

Last time,

defined  $\mathcal{M}G_n$  = moduli space of graphs  $G$   
 w/  $G$  **admissible** := **connected**,  $|V| \geq 3$ ,  $\chi = 1-n$

$$= CV_n / \text{out } F_n$$

also  $CV_n^*$ ,  $\mathcal{M}G_n^* = CV_n^* / \text{out } F_n$

$$\partial CV_n^* = CV_n^* \setminus CV_n$$

$$\partial \mathcal{M}G_n^* = \mathcal{M}G_n^* \setminus \mathcal{M}G_n$$

We observed

$$CG_* = \bigoplus_{n \geq 2} CG_*^{(n)} \quad (\chi = 1-n)$$

and identified

$$C_*(\mathcal{M}G_n^*, \partial \mathcal{M}G_n^*) \leftrightarrow CG_*^{(n)} \quad (\theta = \text{Comm, even or})$$

grading here is by #edges - 1      grading here is by #vertices

If  $G$  has  $k+1$  edges and  $\chi = 1-n$ , then

$G$  has  $2+k-n$  vertices

so

$$C_k(\mathcal{M}G_n^*, \partial \mathcal{M}G_n^*) \leftrightarrow CG_{2+k-n}^{(n)}$$

on both sides, a generator  $G$  is zero if it has an odd automorphism.

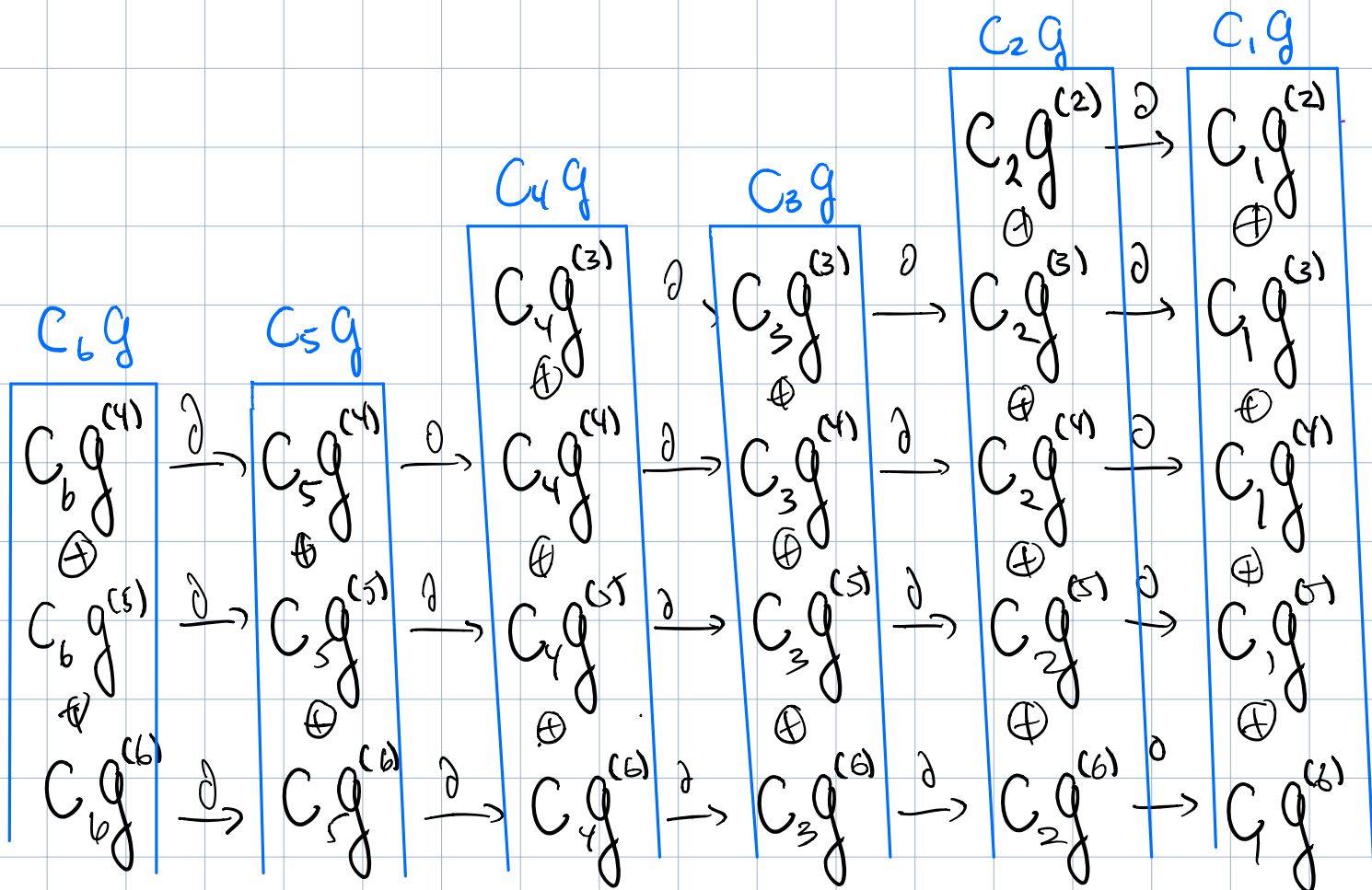
Since every graph in  $CV_n$  has  $\geq n$  edges, the entire  $(n-2)$ -skeleton of  $CV_n^*$  is contained in  $\partial CV_n^*$ .

$$\text{So } C_k(MG_n^*, \partial MG_n^*) = 0 \text{ if } k \leq n-2$$

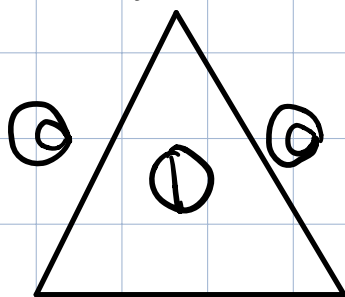
$$\parallel \\ Cg_{2+k-n}^{(n)} = 0 \text{ if } k \leq n-2$$

$$\text{ie } Cg_0 = Cg_{-1} = Cg_{-2} = 0.$$

The entire chain complex  $C_*g$  looks like

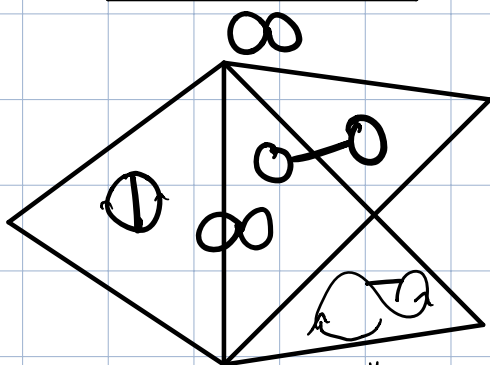


$\partial$  operator collapses edges:



2 vertices become one vertex, lose an edge

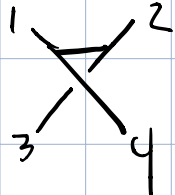
$\delta$  operator splits vertices:



one vertex splits into 2 vertices, add an edge

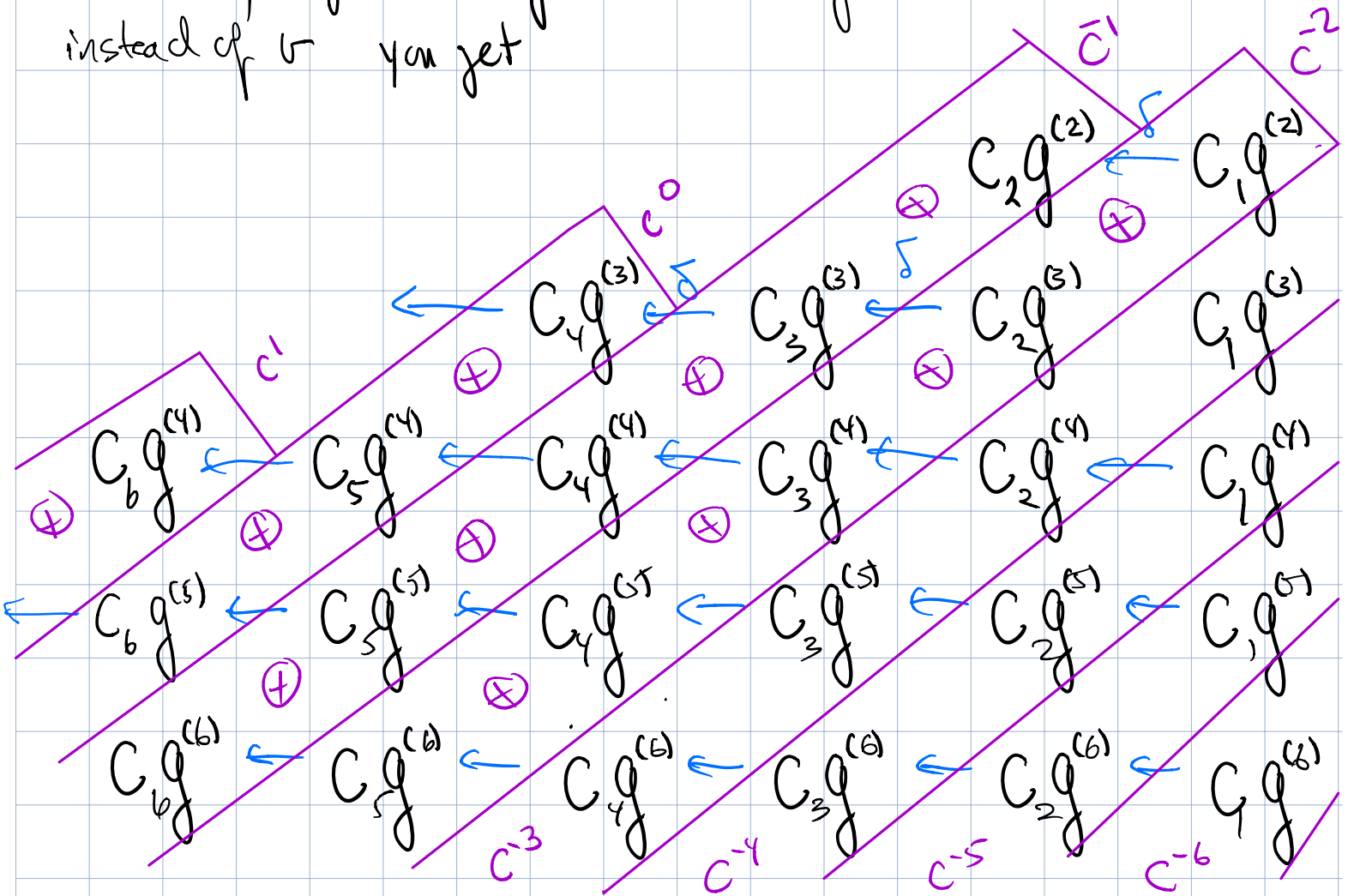
eg  $\delta X_4^2$

$$= \begin{matrix} 1 \\ \diagdown \quad \diagup \\ 3 \quad 4 \end{matrix}^2 + \begin{matrix} 1 \\ \diagup \quad \diagdown \\ 3 \quad 4 \end{matrix}^2$$



both  $\partial$  and  $\delta$  preserve  $\chi = 1 - n$

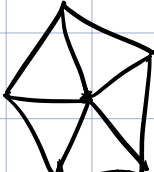
If you grade  $Cg^*$  so that  $\deg G = v - (n+1)$  instead of  $v$  you get



You get Willwacher's graded (co) chain complex

$$H^0 = C_{n+1} g^{(n)} / \text{im}(C_n g^{(n)})$$

$$C_0 = \ker (C_{n+1}^{(n)} \xrightarrow{\partial} C_n^{(n)})$$

contains  $W_n =$  

even  $\Rightarrow (W_n, \sigma) = 0$

odd  $\Rightarrow (W_n, \sigma) \neq 0$



Willwacher found cocycles  $\sigma_n \in C^0, n \text{ odd}$

(he identified  $H^0$  with  $\text{grt}_1$ )

= "Grothendieck-Teichmüller Lie algebra (unipotent version)"

F. Brown proved  $\text{grt}_1$  contains a free Lie algebra generated by odd classes

$\sigma_3, \sigma_5, \dots$

Willwacher showed  $\sigma_n(\omega_n) \neq 0$  ( $n \text{ odd}$ )

If  $G$  has  $k$  vertices and rank  $n$ , it has  $k+n-1$  edges  
so  $\dim \mathcal{S}(G, g) = k+n-2$

Willwacher's result has  $k=n+1$ , so translates to

$$H^{2n-1}(\mathcal{M}g_n^*, \partial \mathcal{M}g_n^*) \neq 0 \quad (n \geq 3)$$

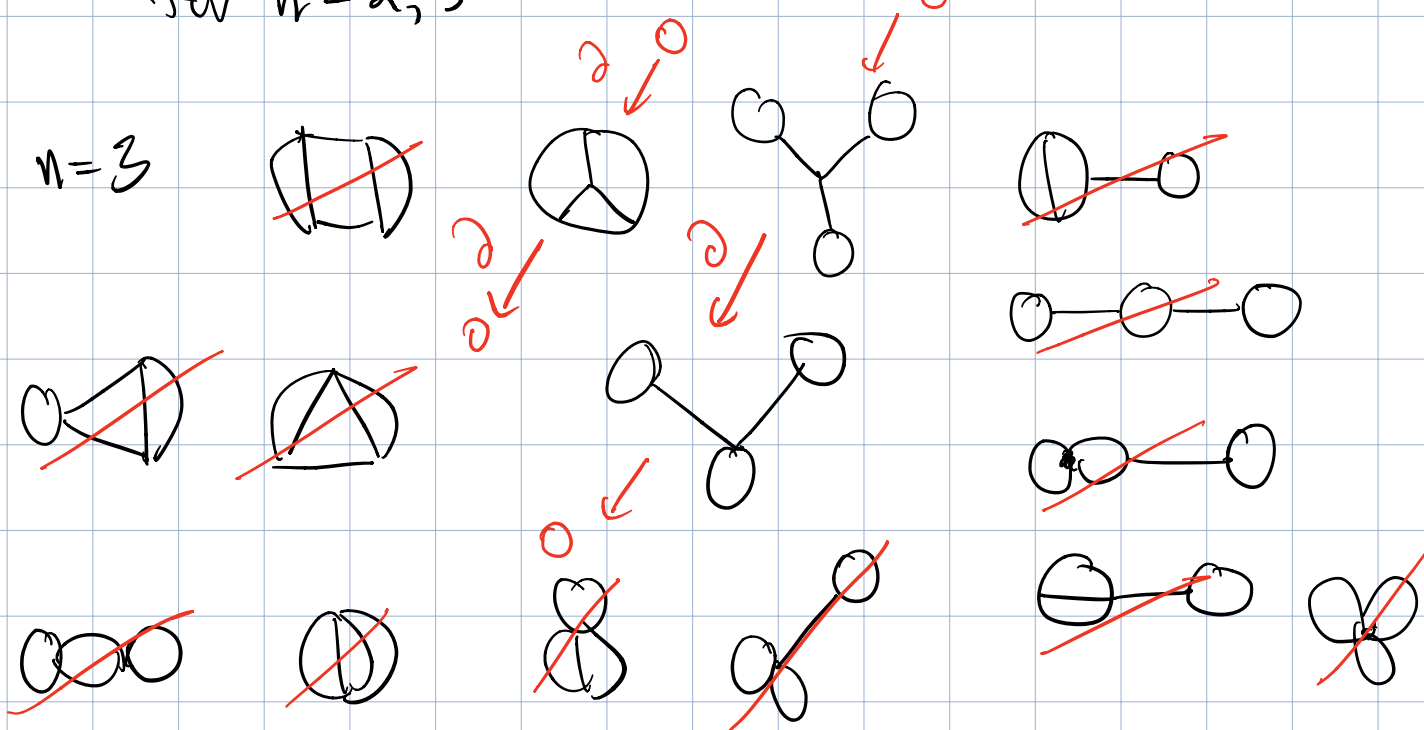
Willwacher also showed  $H^k(Cg_n^*) = 0$  for  $k < 0$

This means  $H^i(\mathcal{M}g_n^*, \partial \mathcal{M}g_n^*) = 0$  for  $i < 2n-1$

eg for  $n=3$ ,  $\dim \mathcal{M}g_3^* = 5$ ,  $H^i = 0$  for  $i < 5$   
so only cohomology is in  $\dim 5$ .

All non-zero homology lies in  $H_k$ ,  $2n-1 \leq k \leq 3n-4$

Exercise: Compute  $H_*(\mathcal{M}g^*, \partial \mathcal{M}g^*)$   
for  $n=2, 3$

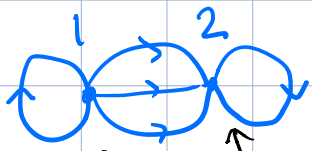


# Back to Kontsevich's theorem:

$$\mathcal{O} = \text{Lie} \Rightarrow H_d(\text{CG}_*^{\text{Lie}}) \cong H^{2n-2-d}(\text{Ort } F_n)$$

where the orientation on  $\mathcal{O}$ -graphs is the odd orientation

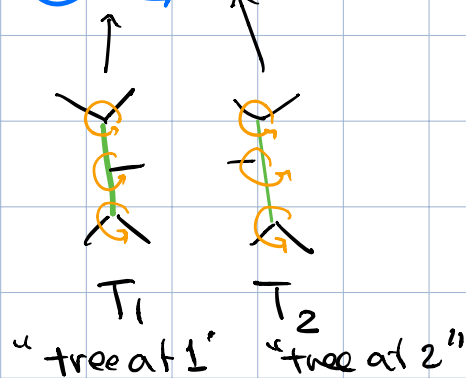
A generator of  $\text{CG}_k$  is a complicated object:

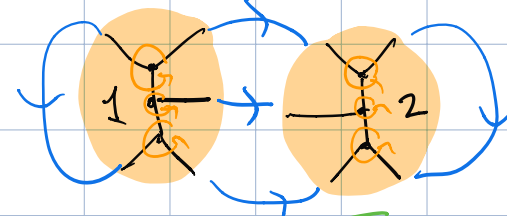
An odd-oriented graph  $X =$  , with

vertices decorated with Lie trees

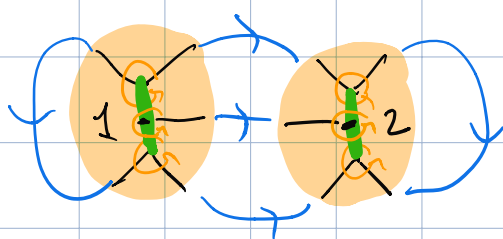
These are Planar binary trees

modulo AS, IHX on internal edges



$$(X, \{T_i\}, \alpha) =$$


Inside each generator is a natural forest  $\Phi$  (= union of trees)



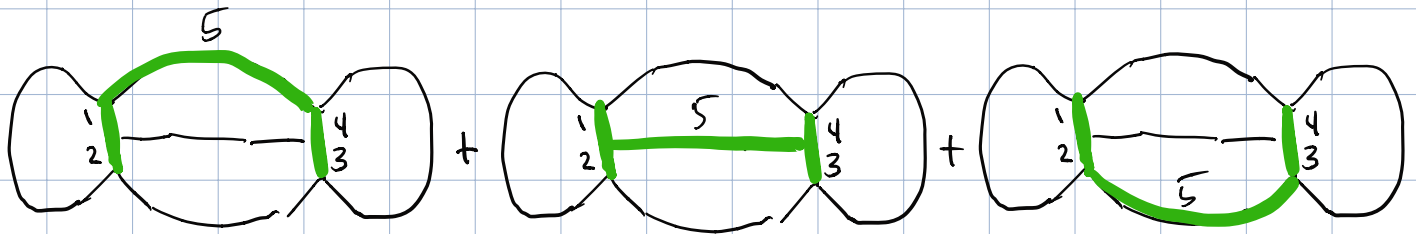
containing all vertices.  
 $\Phi =$  union of internal edges of the trees  $T_i$

Theorem (Conant-V): All of this orientation data is equivalent to ordering the edges of the forest

$$(G, \underline{\Phi}, \alpha) = \text{Diagram} \quad \mathbb{I}$$

Furthermore, the  $\partial$  operator sums over adding an edge to the forest (labeled by next number)

$$\partial(G, \underline{\Phi}, \alpha) =$$



Exercise  $\partial^2 = 0$

Proof of Theorem is linear algebra

(I don't know a direct way of seeing this)

Recall an orientation determined by an ordering of  $S = \{x_1, \dots, x_n\}$  can be described as a choice of unit vector in  $\wedge^n \mathbb{R}S$ , where  $\mathbb{R}S$  is the  $\mathbb{R}$ -space of basis  $S$ .

(ordering for  $x_i$  gives  $x_1 \wedge \dots \wedge x_n \in \wedge^n \mathbb{R}S \cong \mathbb{R}$   
 $:= \det \mathbb{R}S$ )

so an orientation on an  $\mathcal{O}$ -graph is a unit vector in

$$\det(\mathbb{R}V(X)) \otimes \bigotimes_{e \in E(X)} \det \mathbb{R}H(e) \otimes \bigotimes_{v \in V(X)} \left( \bigotimes_{u \in V(T_v)} \det \mathbb{R}H(u) \right)$$

$\uparrow$   
 order  $v(x)$

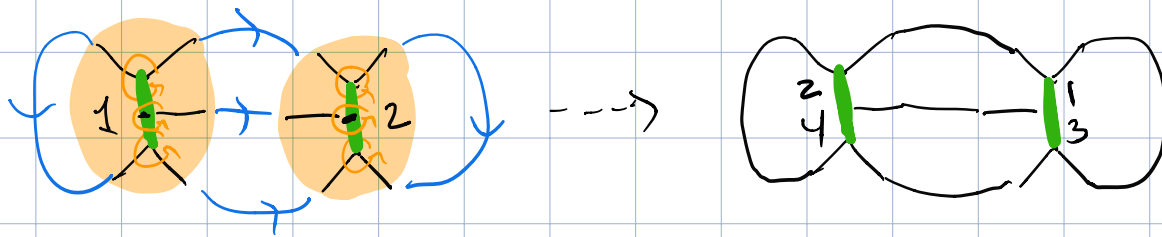
$\uparrow$   
 order each pair  $H(e)$  of half-edges for  $e \in X$

$\uparrow$   
 order the half-edges at each vertex of each tree  $T_v$

We are claiming there is a canonical isomorphism of the above expression with

$$\det \mathbb{R}E(\Phi) \quad (E(\Phi) = \text{edges of } \Phi)$$

(ie picking a unit vector in either side determines a unit vector in the other side)



I don't know a direct way to see this  
 Instead we use two lemmas

**Lemma 1**:  $S = \coprod_{i=1}^k S_i$  a finite set

Then  $\exists$  a canonical isomorphism

$$\bigotimes_i \det S_i \cong \det \left( \bigoplus_{|S_i| \text{ odd}} \mathbb{R} \right) \xrightarrow{\cong} \det \mathbb{R}^S$$

$\uparrow$  order of each  $S_i$        $\uparrow$  order the set of  $S_i$  with  $|S_i|$  odd       $\rightarrow$  total ordering of  $S$

Pf switching  $x, y \in S_i$  or  $S_i \leftrightarrow S_j$  with  $|S_i|, |S_j|$  odd changes the sign of the total ordering of  $S$

$S_i \leftrightarrow S_j$  doesn't change the sign if  $|S_i|$  or  $|S_j|$  is even.

(map to other way "groups" the elements of  $S$ )

Lemma 2  $0 \rightarrow U \xrightarrow{f} V \xrightarrow{p} W \rightarrow 1$

a short exact sequence of finite-dim'd vector spaces. Then  $\det V$  is canonically isomorphic to  $\det U \otimes \det W$ .

Pf

Choose a splitting  $V \begin{matrix} \xrightarrow{\circ} \\ \xrightarrow{p} \end{matrix} W$

The isomorphism is given by

$$\begin{array}{ccc} \det U \otimes \det W & \longrightarrow & \det V \\ u \otimes w & \longmapsto & f(u) \wedge \Delta(w) \end{array}$$

This is independent of  $\circ$ , since  $ps = \text{id}$

Exercise  $0 \rightarrow U \rightarrow V \rightarrow W \rightarrow Z \rightarrow 0$

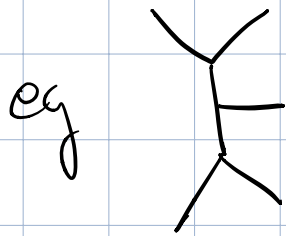
a short exact seq of finite-dim'd v. spaces

$\Rightarrow$   $\exists$  canonical isomorphism

$$\det U \otimes \det W \cong \det V \otimes \det Z$$

Hint: split the sequence into two short exact sequences)

# How to use these lemmas:



claim: cyclic orderings around each vertex

$T$  a binary tree  $\iff$  ordering all edges

(The "obvious" correspondence doesn't work)

Use the augmented chain complex of  $T$ :

$$0 \rightarrow C_1 \rightarrow C_0 \rightarrow \mathbb{R} \rightarrow 0$$

This is exact ( $\tilde{H}_* T = 0$ )

$$C_0 = \mathbb{R}V \quad \det C_0 = \det \mathbb{R}V$$

To give a chain in  $C_1$ , you need to prescribe orientations on the edges, ( $\partial e = t(e) - i(e)$ )

$$\det C_1 = \det \mathbb{R}E \otimes \bigotimes_e \det H(e)$$

Now use Lemma 1:

$$\det \mathbb{R}V \cong \det \mathbb{R}E \otimes \bigotimes_e \det H(e) \otimes \cancel{\det \mathbb{R}}$$

only one ordering of a one-element set!



$$\text{and Lemma 2} \quad \cong \det RE \otimes \det RH$$

Tensor both sides w/  $\det RV$ ,  $\det RE$

$$\det RE \cong \det RV \otimes \det RH$$

Now use Lemma 2 again

$$\det RE \cong \det RV \otimes \left( \bigotimes_{v} \det H(v) \otimes \det \left( \bigoplus_{|v| \text{ odd}} \mathbb{R} \right) \right)$$

all vertices are odd  $\rightarrow$   $|v| \text{ odd}$

$$\cong \det RV \otimes \left( \bigotimes_{v} \det H(v) \right) \otimes \det RV$$

$$\cong \bigotimes_{v} \det H(v) \quad \checkmark$$

The rest of the proof that

$$(X, \{T_v\}, \sigma_X) \sim (G, \Phi, \sigma_\Phi)$$

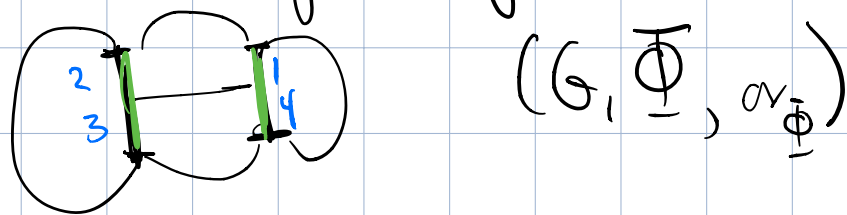
is similar.

(Note that the half-edges of  $X$  are exactly the leaves of the  $T_v$ .)

Also have to check:

$$\partial(X, \{\tau_x\}, \sigma_X) = \sum_{\substack{\Phi \subseteq e \\ \text{a forest}}} (G, \Phi, \sigma_{\Phi})$$

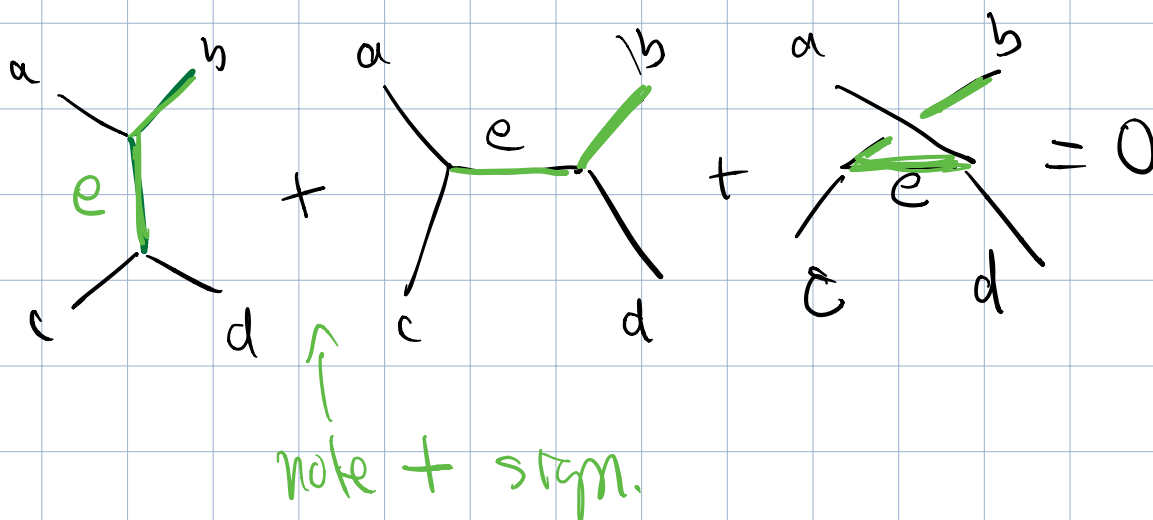
A generator of  $CG_X$  is now



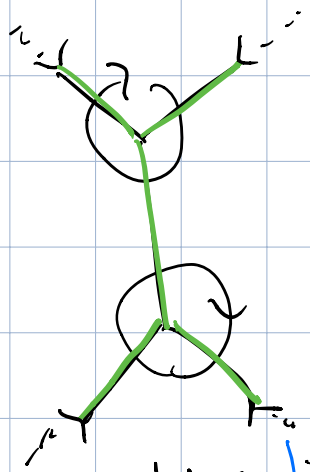
$\Phi =$  forest cutting all vertices of  $G$

$$\sigma \in \det RE(\Phi)$$

IHX translates into

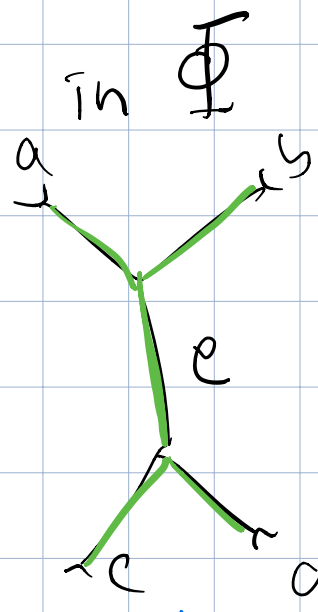


in Lie tree



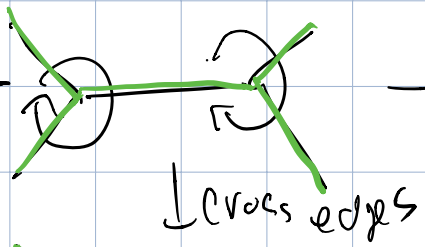
$\downarrow$

turn left

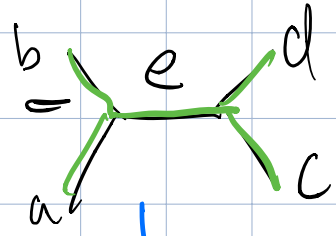


$\downarrow$

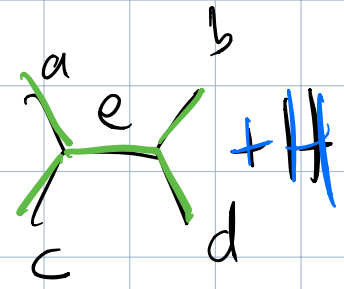
$-H$



cross edges

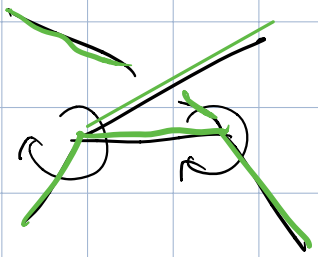


$=$

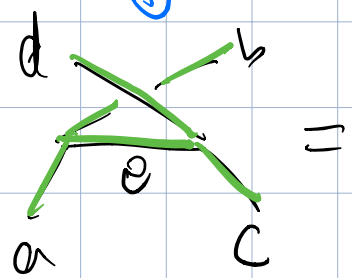


$+H$

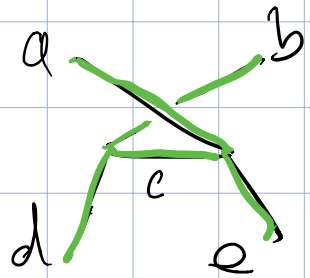
$+X$



$\rightarrow$



$=$



$+X$

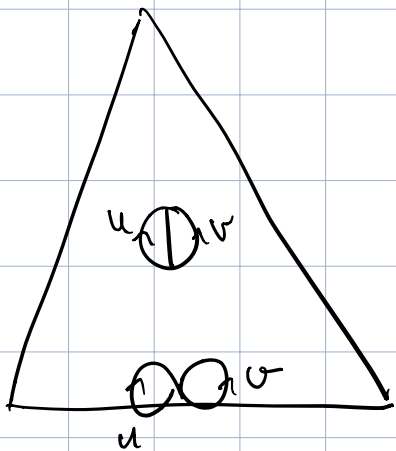
## Back to Outer space $CV_n$

To get the isomorphism

$$\begin{aligned} H_d(\mathbb{C}G_{*}^{(n)}) &\cong H^{2n-2-d}(\text{Out } F_n) \cong H^{2n-2-d}(\mathbb{M}g_n) \\ &= H^{2n-2-d}(CV_n / \text{out } F_n) \end{aligned}$$

It's easiest to look at  $CV_n$  first:

$CV_n =$  disjoint union of open simplices  $\sigma(G, g)$ ,  $G =$  admissible, rank  $n$   
(connected,  $|G| \geq 3 + \nu$ )

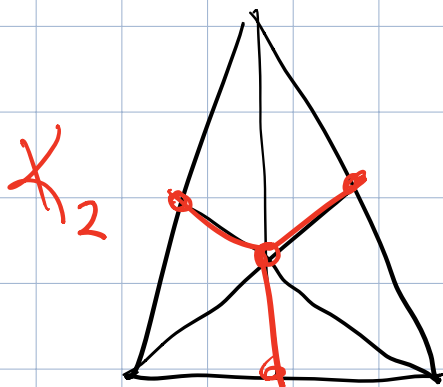


some faces are missing.

$CV_n \subset CV_n^* =$  simplicial completion

$\{u, v\} =$  basis for  $F_n$

Def  $K_n \subset (CV_n^*)'$  (barycentric subdivision). Vertex = simplex of  $CV_n^*$ .



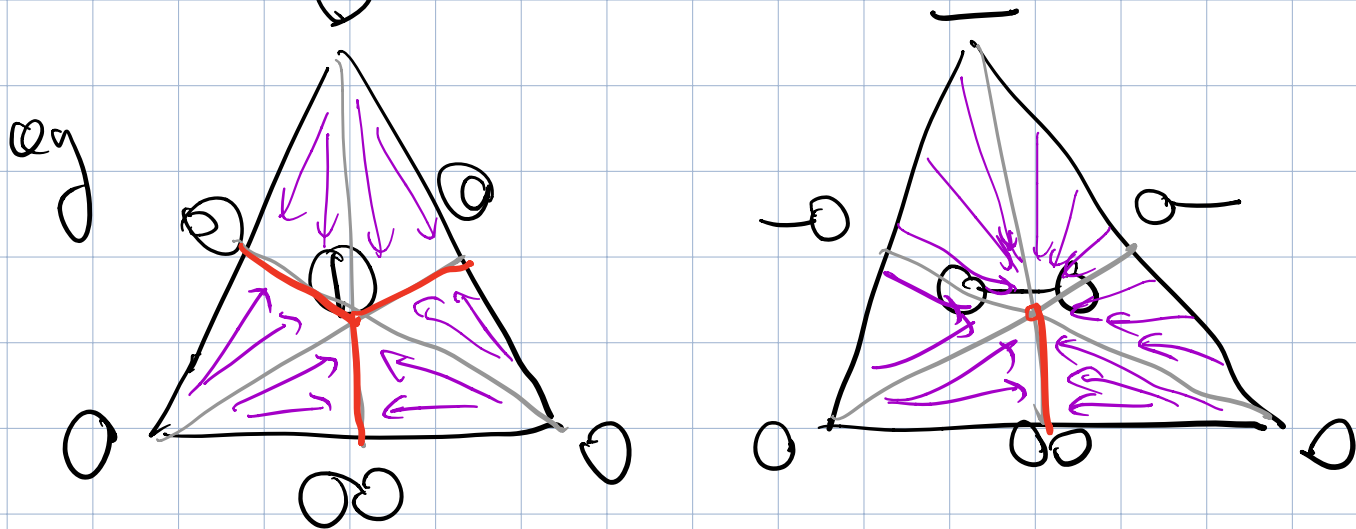
= span of vertices in  $CV_n$  ( $\leftarrow$  open simplices in  $CV_n$ )

= geometric realization of poset of  $\sigma(G, g)$

All maximal simplices of  $(CV_n^*)'$  have a missing face (at least a vertex) and a face in  $K_n$  (at least a trivalent graph)

$CV_n \searrow K_n$  (deformable retract)

by linearly retracting each maximal simplex of  $(CV_n^*)'$  to its face in  $K_n$



The action of  $\text{Cut}(F_n)$  permutes the  $\sigma(G, g)$ ,  
 so  $K_n$  is invariant under the action

$K_n$  is a def-retract of  $CV_n$ , so  
 is contractible

$$\Rightarrow H^*(\text{Cut } F_n) = H^*(K_n / \text{Cut } F_n)$$

$K_n$  is a simplicial complex  
 maximal simplex = chain of edge-collapses.

( $n=2$  picture is not enlightening here!  
 Try an  $n=3$  picture)