

We've shown the homology of Kontsevich's Lie graph complex (odd version) $CG_{\neq}^{(n)}$ is equal to the cohomology of $K_n / \text{Out } F_n$

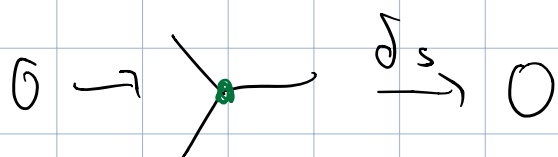
(though the chain complexes are not the same, and in fact look very different at first)

We left out one step.. showing the vertical columns in the double complex have no homology except in the top dimension.

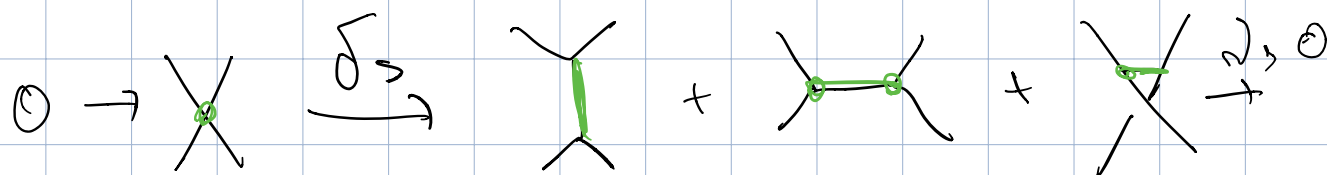
The vertical maps δ_s split one vertex-into two, then add. What happens at one vertex is independent of what happens at another.

So let's just look what happens at one vertex v in one graph G .

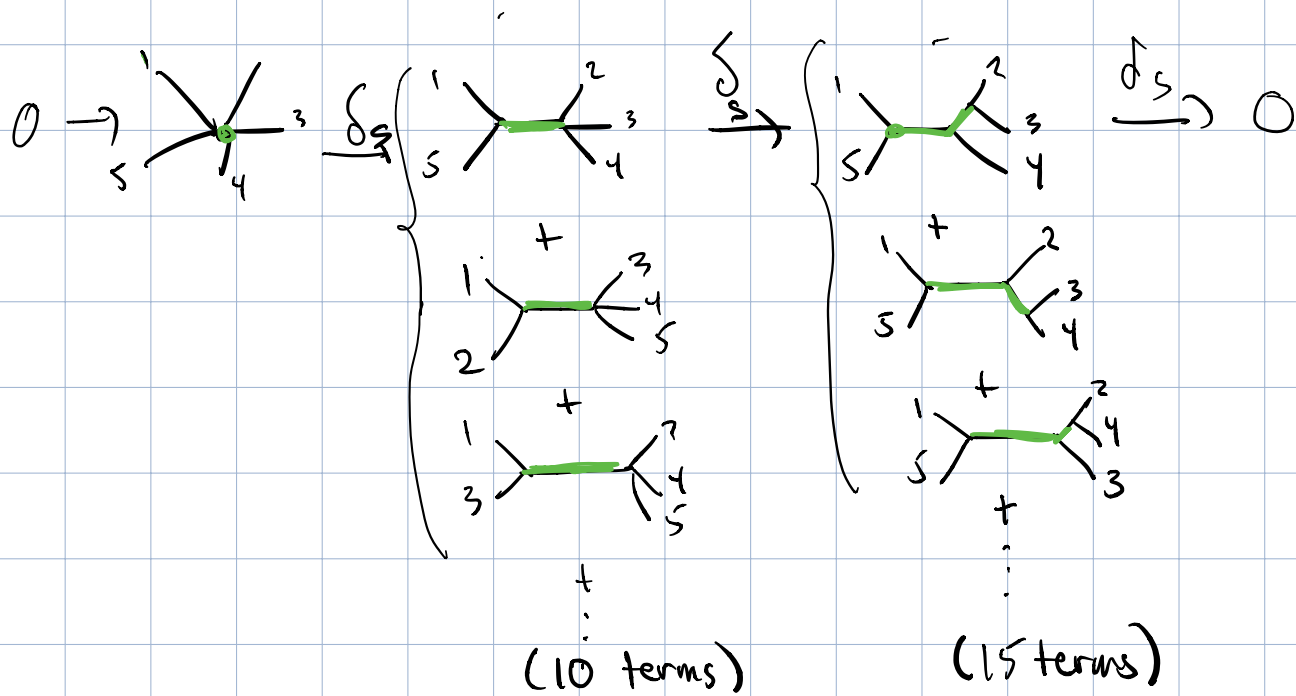
If $|v| = 3$, δ_s doesn't affect it.



If $|v| = 4$



If $|V| = 5$, you can split twice:



$$(*) \quad 0 \rightarrow \mathbb{R} \xrightarrow{\delta_0} \mathbb{R}^{10} \xrightarrow{\delta_1} \mathbb{R}^{15} \rightarrow 0$$

(*) is the augmented cochain complex of a simplicial complex T_n with

- a 0-cell for each 1-edge tree with n labeled leaves

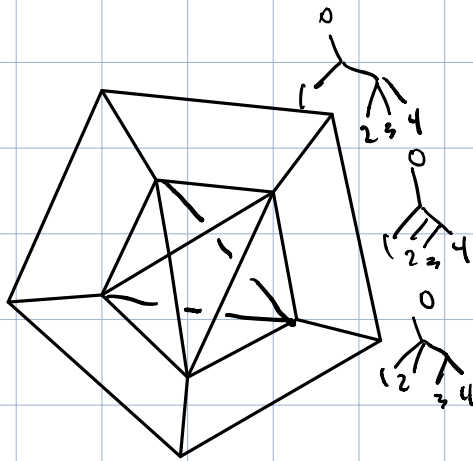
- a 1-cell for each 2-edge tree with n labeled leaves

- an $(n-4)$ cell for each $(n-3)$ -edge tree with n labeled leaves.

of
with $n=4$, this s.c has 3 0-cells

$$\cong S^0 \vee S^0$$

with $n=5$ we have 10 0-cells
15 1-cells,



looks like
(Petersen graph)

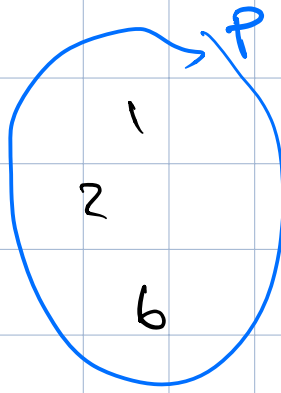
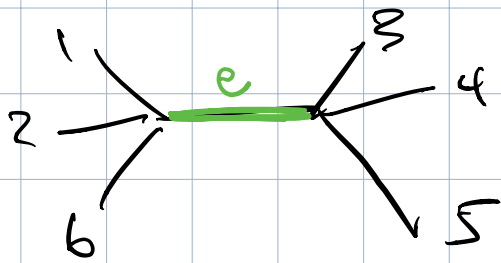
Prop $T_n \cong \vee S^{n-4}$ for any n .
(ie has $H_k = 0$ unless $k = n-4$)

PF

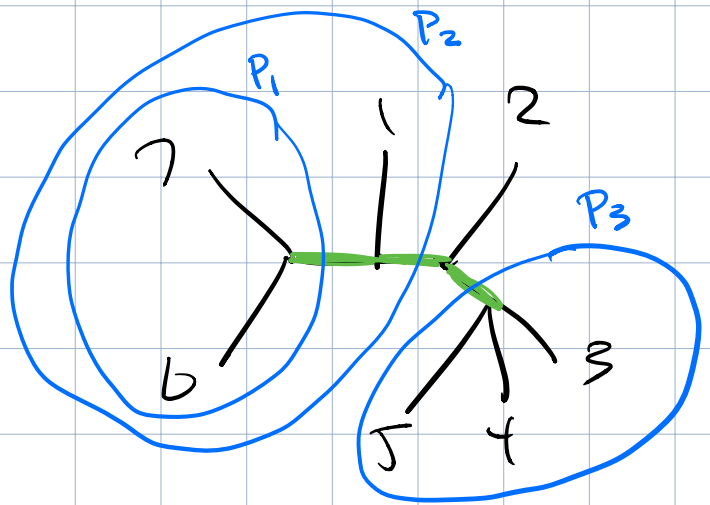
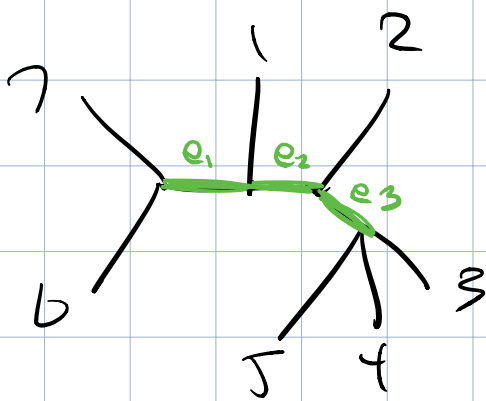
Another way to describe a tree
with n labeled leaves and 1 edge is
as a **thick partition** of the
set $\{1, \dots, n\}$:

(thick: each side has ≥ 2 elements)

eg



With k edges, you get k pairwise compatible partitions:



(P, Q compatible \Leftrightarrow some side of P is disjoint from some side of Q)

\Leftrightarrow partitions can be drawn in the plane so that the circles don't intersect.

Now let's reconsider the simplicial complex T_n :

- vertex for each partition (= 1-edge tree)
- edge for each pair (P, Q) of compatible partitions

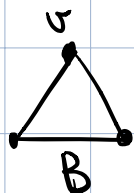
T_n is flag: $\{P_0, \dots, P_k\}$ are pairwise compatible if and only if they span a k -simplex

Exercise let $B \subset X$ be a full subcomplex of a flag complex.

If $v \in X \setminus B$ is a vertex with $\text{ll}v \cap B \neq \emptyset$, and $J \subset B$ is the subcomplex spanned by $\text{ll}v \cap B$ then the subcomplex $\langle B, v \rangle$ spanned by B and v is equal to $B \cup_J c(J)$.

(non-examples: $X = \partial$ of a triangle, (not flag)

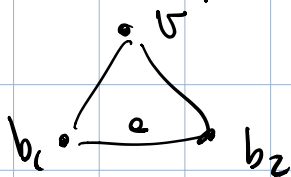
$B = \text{edge}$.



Then $J = B$

but $X \neq B \cup_B c(B)$

$X = \text{triangle}$, $B = 2 \text{ vertices (not full)}$



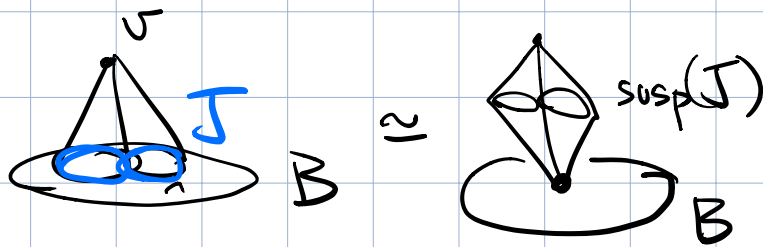
$J = \{b_1, b_2\}$ $cJ = \text{edge}$

$B \cup_J c(J) = \partial X$

Corollary B, X, ν , as above. If B is contractible then $\langle \nu, B \rangle \cong \text{susp } J$
 If $B \cong VS^k$ and $J \cong VS^{k-1}$ then $\langle \nu, B \rangle \cong VS^k$.

pf Mayer-Vietoris plus van Kampen

picture:

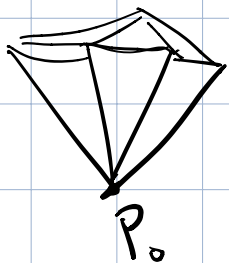


Proof of theorem ($T_n \cong VS^{n-4}$)

Induction on n (true for $n=4, 5$)

Let $P_0 = \begin{pmatrix} 1 \\ 2 \end{pmatrix} \dots n$ $\left(\rightarrow \begin{array}{c} 3 \\ \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \\ 2 \end{array} \right)$

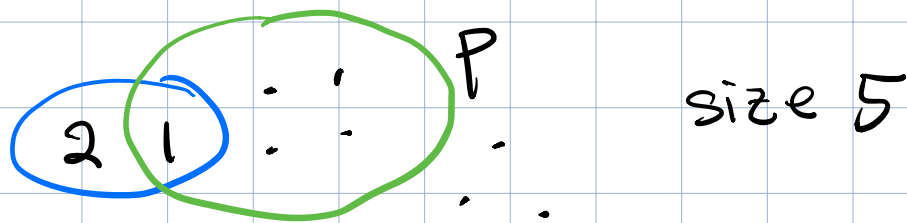
$T_n^0 =$ subcomplex spanned by partitions compatible with P_0



$=$ cone on P_0 , so $\cong \text{pt}$.

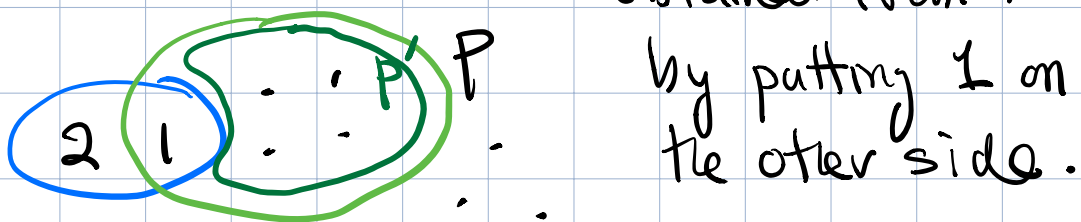
Note $\forall P_0 \cong T_{n-1}$. (shrink the leaves $1:2$ to a pt)

Not in T_n^0 : partitions that cross P_0



Define the size of P to be the number of elements in the side containing 1.

If P crosses P_0 , let P' be the partition obtained from P



P' is compatible with both P and P_0

We will add all P of size > 2 to $c(P_0)$ in order of decreasing size, using the corollary to keep control of the homotopy type.

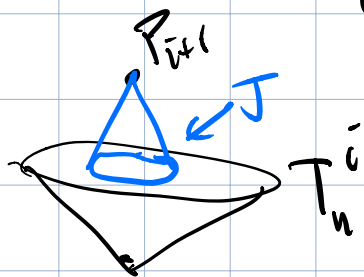
So... Order the $P \in T_n \setminus c(P_0)$ so that

$\text{size}(P_1) \geq \text{size}(P_2) \geq \dots$

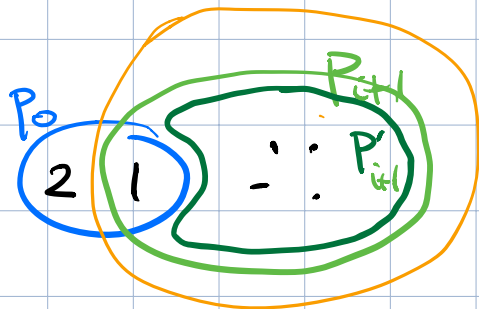
let $T_n^c = \text{span of } T_n^0 \text{ and all } P_j \text{ with } j \leq i$

Claim: $J = \langle \text{all } P_{i+1} \cap T_n^i \rangle$ is a cone on P_{i+1}'
 so is contractible.

PF: $Q \in J$ if Q is
 compatible with P_{i+1}'
 and either

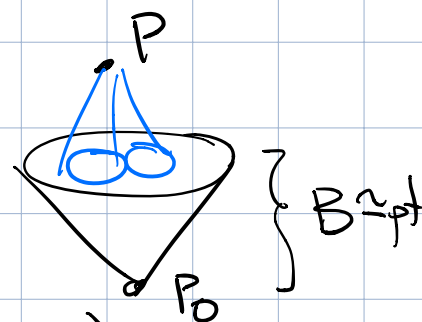
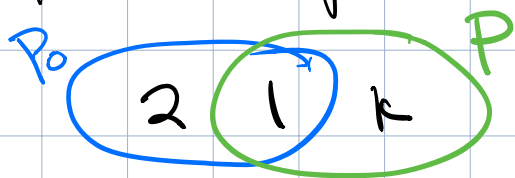


has size \geq size P_{i+1}'
 (which \Rightarrow size $>$ size P_{i+1}')
or is compatible with P_0



In either case, Q is compatible with P_{i+1}'
 so $J = c(P_{i+1}')$

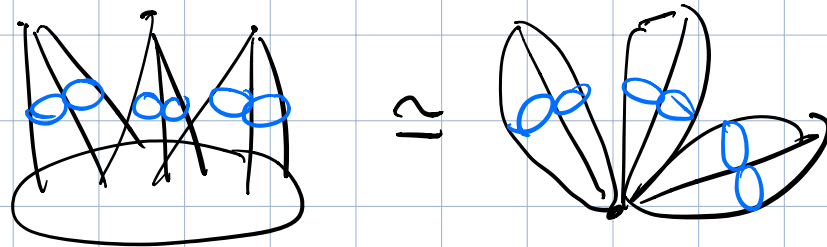
The only vertices we haven't yet included are
 partitions of size 2 that cross P_0 !



But $\text{all } P \cap B = \text{all } P$
 $\cong T_{n-1}$ (as we saw)
 $\cong VS^{n-5}$ by induction

so adding P gives $\text{Susp}(VS^{n-2}) \cong VS^{n-1}$.

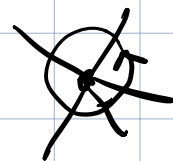
Adding all other $P = \textcircled{1k}$ still gives VS^{n-1} !



Next: We've related commutative and Lie graph complexes to the topology of moduli spaces of graphs.

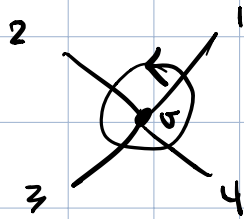
What about associative graphs?

Structure at a vertex = cyclic ordering of adjacent edges



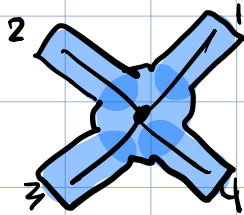
A graph with a cyclic ordering of the edges at each vertex can be "fattened" into a unique oriented compact connected surface

First put a neighborhood of each vertex
into the plane, oriented counterclockwise

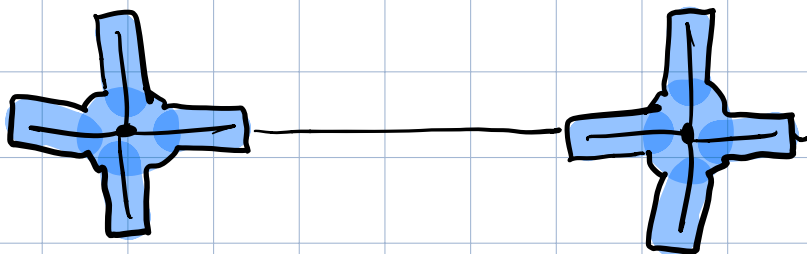


$\subset \mathbb{R}^2$

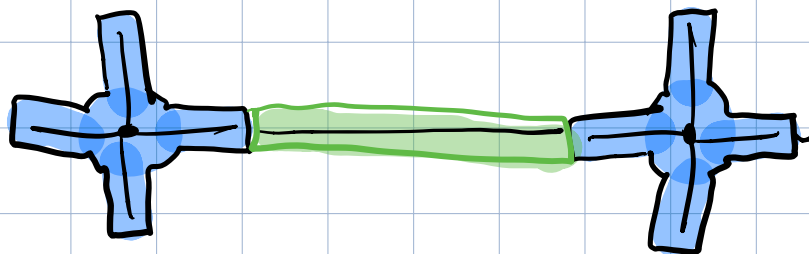
Then fatten this into a disk with tabs

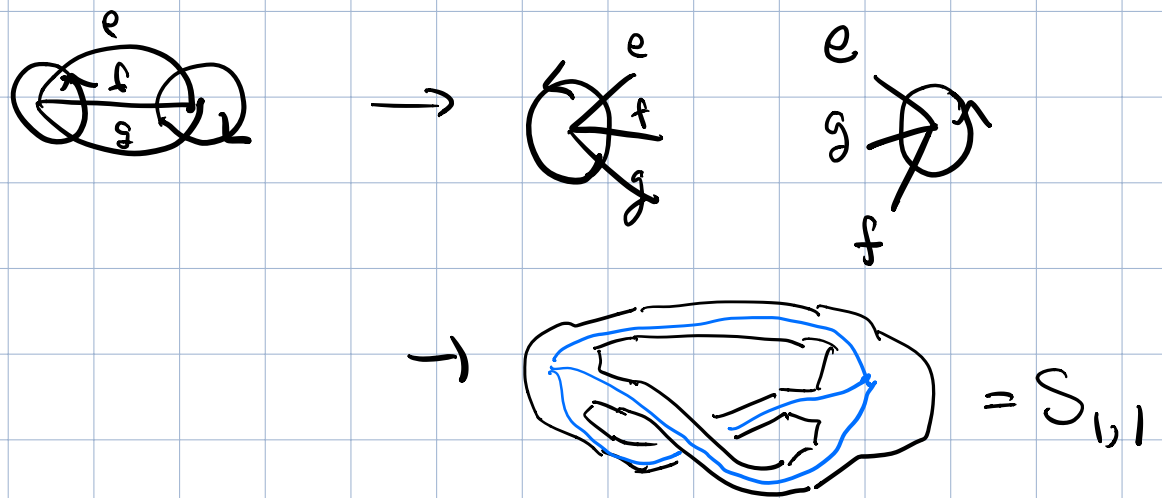
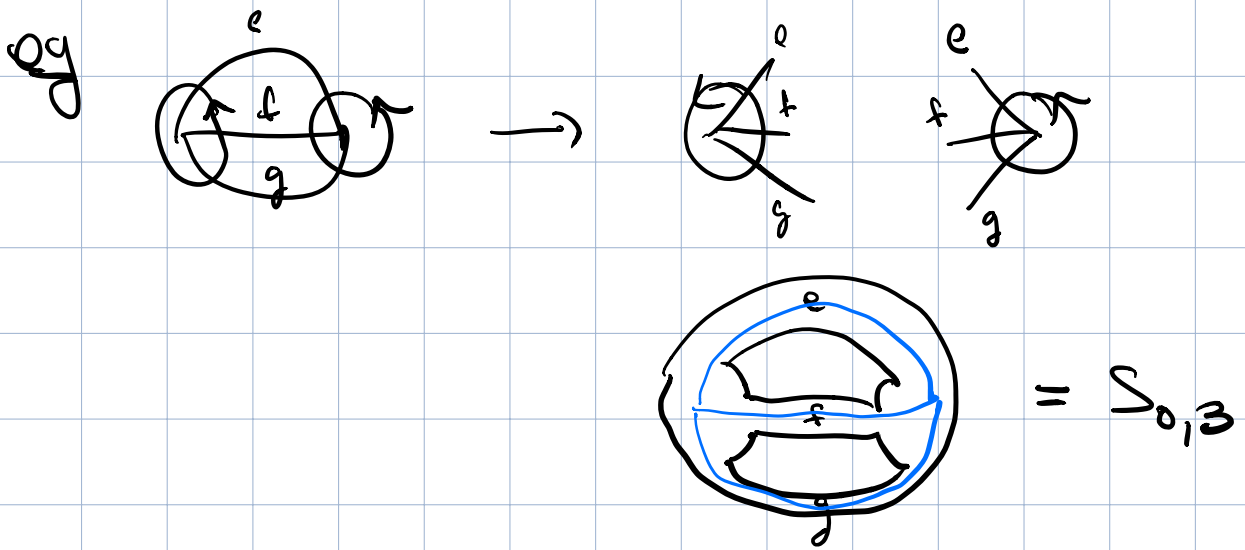


The edges of the graph
give a pairing of the tabs



Connect each pair with an (oriented) rectangle





Collapsing an edge doesn't change the homeomorphism type of the surface, so the Assocative graph complex breaks up into a direct sum

$$\text{Cg}_* = \bigoplus_{S = \text{surface}} \text{Cg}_*^S$$

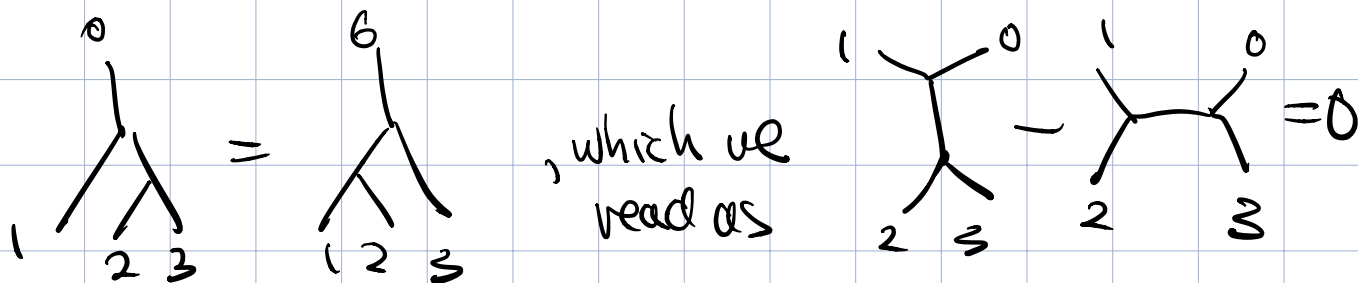
(compact, oriented, connected, with ∂)

We've already done almost all the work needed to relate CG_* to a moduli space of graphs, when we studied Lie graphs

We described generators of the Lie operad as planar trivalent trees modulo AS and IHX relations

We described generators of the Associative operad as planar "stars" ~~IHX~~

But back up... this is equivalent to planar trivalent trees modulo the associative relation

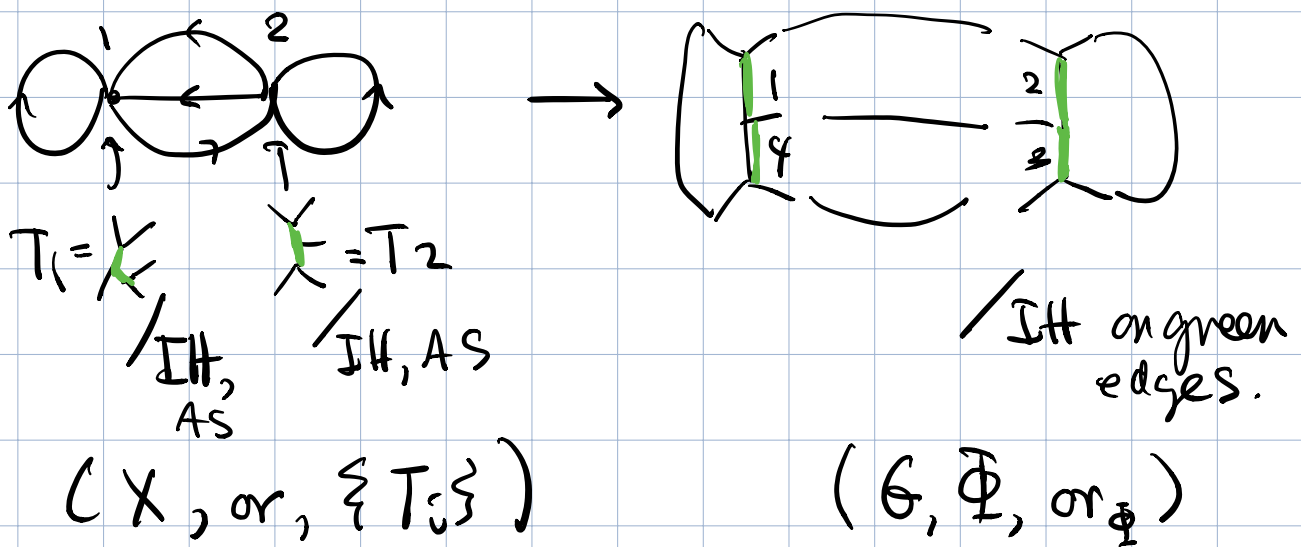


"IHX"-relation"

A cyclic ordering of the edges in a star is clearly \sim to a cyclic ordering of the leaves of this trivalent tree

We saw this is \sim ordering the edges of the tree

An associative structure on an odd-oriented graph is



Associative graph \longrightarrow forested graph

We need to check the orientation lemma

Lemma: The orientation data on $(X, \text{or}, \{T_i\})$ graph is equivalent to an ordering of the edges of Φ in (G, Φ, α) , and the ∂ operator is the same

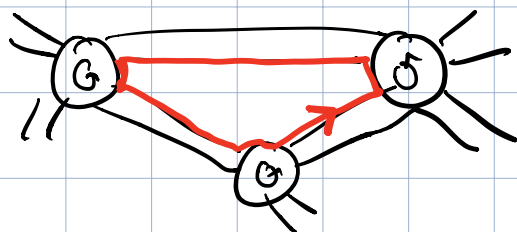
(We didn't completely prove this in the Lie case - I just gave you the tools. Ditto here.)

This shows Cg_*^S is isomorphic to $f g_*^S$,
 so they have the same homology

We now claim $f g_*^S$ has the same
 homology as the cochain complex
 $C^*(K^S / \Gamma(S))$, where $K^S =$ graphs that
 fatten to S . We have

$K^S \hookrightarrow K_n$ ($\subseteq CV_n$, $\pi_1 S \cong F_n$) by
 forgetting the cyclic order at each
 vertex of a ribbon graph.
 and

$\Gamma(S) \hookrightarrow \text{Out}(F_n)$ is the
 subgroup that "preserves the cyclic order
 at each vertex", i.e.
 $\Gamma(S)$ preserves the "boundary words",
 the cyclic words in F_n corresponding
 to the boundary curves in S :



To prove the claim, we decompose the δ operator in $C^*(K^S)$ as before,

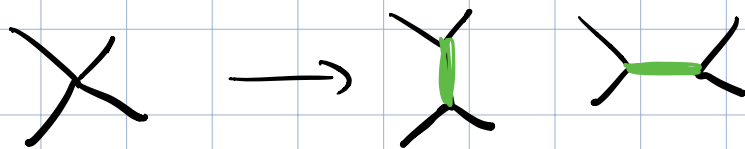
$$\delta = \delta_a + \delta_s \quad \begin{array}{l} \delta_a \text{ adds an edge to } \emptyset \\ \delta_s \text{ splits a vertex of } G \end{array}$$

and, as before, we show the vertical subcomplexes (using δ_s) have homology only in the top dimension

Exercise Identify the vertical (co)chain complex as the cochain complex of a sphere:

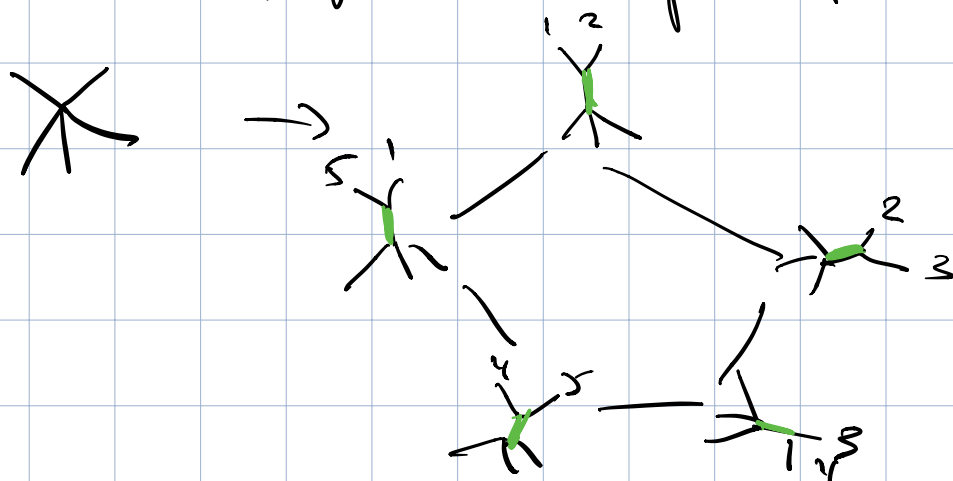
look at the trees you can get by IH splitting, starting at a single vertex v of valence $|v|$.

If $|v| = 4$, there are now only two possible splittings without leaving S :



so cplx is $0 \rightarrow \mathbb{R} \rightarrow \mathbb{R}^2 \rightarrow 0$
= augmented cochain cplx of S^0

If $|S|=5$, get the ∂ of a pentagon



Now

(1) Zieschang proved $\Gamma(S) \cong \text{Mod}(S)$

(2) The proof that K_n is contractible restricts without change to K_S

$\therefore K_S / \Gamma(S)$ is a torus

$K(\text{Mod}(S), 1)$,

$$\therefore H_d(CG_x^S) = H_d(fG_x^S) = H^{2r-2-d}(\text{Mod}(S))$$