## Summer School - Number Theory for Cryptography

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- 1. a) Program the p-1 method, together with its standard continuation.
  - b) Program Brent's continuation.
- c) Let us describe a Baby-Step-Giant-Step continuation. Select  $w \approx \sqrt{B_2}$ ,  $v_1 = \lceil B_1/w \rceil$ ,  $v_2 = \lceil B_2/w \rceil$ . Write our prime s as s = vw u, with  $0 \le u < w$ ,  $v_1 \le v \le v_2$ . Using the fact that  $\gcd(b^s 1, N) > 1$  iff  $\gcd(b^{wv} b^u, N) > 1$ , give an algorithm to perform the search for s in  $O(\sqrt{B_2})$  operations. Implement it and compare it to the continuations already given.
- 2. Let p = 2k + 1 and q = 4k + 1 be prime (with  $k \ge 5$ ) and N = pq. Show that for all a prime to N, one has  $\gcd(a^{k!} 1, N) = N$ . Show how to modify the p 1 method of factoring to recover p and q in that special case.
- 3. Give parametrizations of elliptic curves over  $\mathbb{Q}$  having a 2-torsion point, respectively torsion subgroup of order 4, in Weierstrass or Edwards form. Are they interesting to use in ECM?
- 4. Program ECM.
- 5. (a) Let  $\omega = 2^s t$ , with integers s and t, t odd. One considers the sequence  $y_0 = 1$ ,  $y_k = 2y_{k-1} \mod \omega$ . Let e be the order of 2 in  $(\mathbb{Z}/t\mathbb{Z})^*$ . Show that the length and tail of the iteration are respectively s and e.
- (b) Let p be a prime and  $f: \mathbb{Z}/p\mathbb{Z} \longrightarrow \mathbb{Z}/p\mathbb{Z}$  such that  $f(x) = x^2$ . Define  $(a_k)$  by  $a_0 = a$  (where  $a \not\equiv 0 \mod p$ ) and  $a_k = f(a_{k-1}) = a^{2^k} \mod p$ . Let  $\omega = 2^s t$  be the order of a modulo p. Compute the cycle length and tail length of the cycle of the sequence  $(a_k)$ . Numerical values: p = 59, a = 2. Is the function f suitable for the  $\rho$  method?
- (c) Let K be a field, and u, v two non-zero elements of K. Show that if u + 1/u = v + 1/v, then u = v or u = 1/v.
- (d) Let  $a \neq 1$ ,  $a \in (\mathbb{Z}/p\mathbb{Z})^*$ . Put  $x_0 = a + 1/a$ ,  $f(x) = x^2 2 \mod p$ ,  $x_k = f(x_{k-1})$ . Hence,  $x_k = a^{2^k} + a^{-2^k}$ . Let  $\omega = 2^s t$  be the ordre of a modulo p. Let (H) denote the assertion: "the equation  $2^\ell \equiv -1 \mod t$  has at least one solution". Prove that if (H) is satisfied, the length of the cycle of  $(x_k)$  is e/2; otherwise, this length is equal to e. Compute the tail length.
- (e) Let  $x_0 \in \mathbb{Z}/p\mathbb{Z}$ . Assume there is no solution to  $x_0 = a + 1/a$  for  $a \in (\mathbb{Z}/p\mathbb{Z})^*$ . Let  $\alpha \in \mathbb{F}_{p^2} \setminus \mathbb{F}_p$  such that  $x_0 = \alpha + 1/\alpha$ . Let  $\omega = 2^s t$  be the ordre of  $\alpha$  in  $\mathbb{F}_{p^2}$  and e the ordre of 2 modulo t. Show that  $\alpha, \alpha^2, \ldots, \alpha^{2^s}$  are all distinct in  $\mathbb{F}_{p^2}$  and that  $\alpha^{2^{s+e}} = \alpha^{2^s}$ . Define  $x_k = f(x_{k-1})$  with  $f(x) = x^2 2 \pmod{p}$ . Show that if (H) is true, the cycle length is e/2 and e otherwise. Compute the tail length. Numerical values: p = 5,  $x_0 = 1$ .
- (f) What happens if one uses  $f(x) = x^2 2$  in the  $\rho$  method?