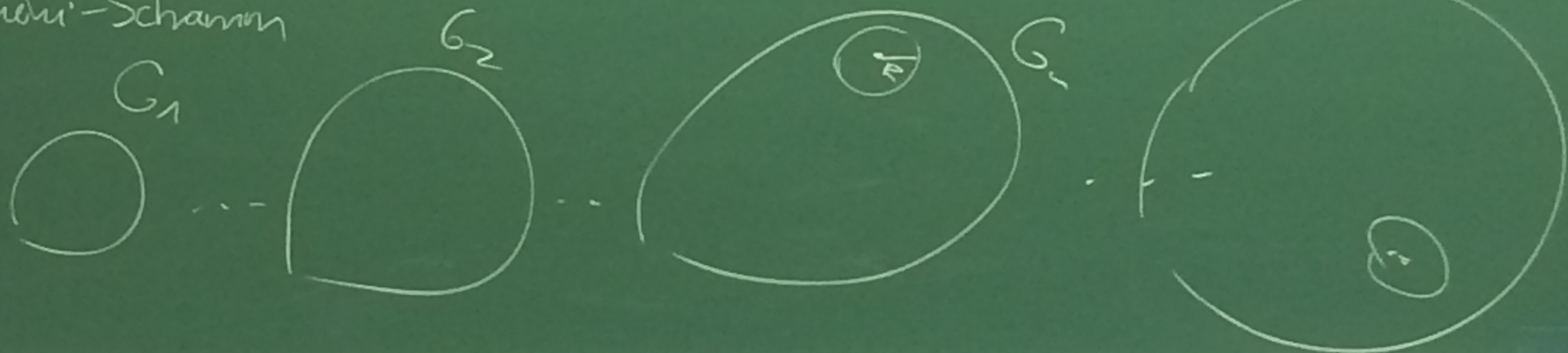


graph : Fix the maximal degree D

G_D : set of graphs \leftarrow FG_D - ~~finite~~ G_D

RG_D : spaces of ^{connected} rooted graphs

Benjamini-Schramm

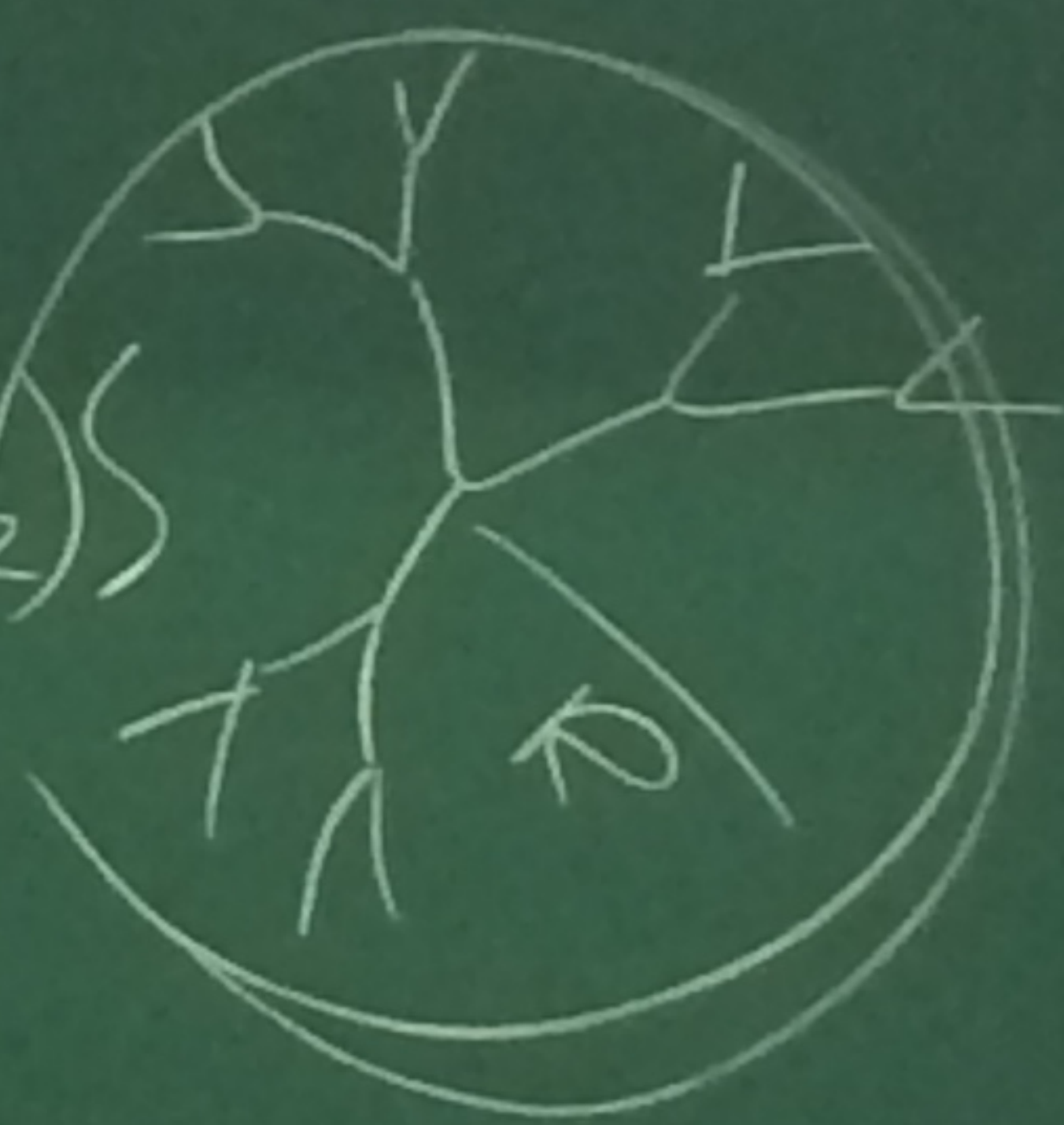


$$G_1, G_2 \in RG_D \quad d(G_1, G_2) = \frac{1}{k} \quad k = \max$$

If G is finite ^{unrooted} graph \Rightarrow

λ_G : random rooted graph

$$\{B_k(G_1) \cong B_k(G_2)\}$$



Def. G_n BS-converges, if λ_{G_n} weakly converges.

Every sequence has a convergent subsequence.

A graph parameter $f: \text{FC}_{\text{FD}}^{\text{finite}} \rightarrow \mathbb{R}$ is testable, if $f(G_n)$ converges, whenever G_n BS-converges.

What is testable? — density of triangles
 — [Lupas] tree entropy: $\frac{\log \# \text{ of spanning trees of } G}{|G|}$

— [Nayyar-Oruk + Eldor-Lippner] $\frac{\text{max size of matching}}{|G|}$ is testable

NON-TESTABLE: independence ratio

LINEAR IND RATIO IS TESTABLE

$\frac{\text{rank}_p(\text{Adj}(G))}{|G|} \leq \mu_G(205)$ is testable

OPEN: mod p !!!

REAL STATEMENT:

μ_G uniform measure on eigenvalues of $\text{Adj}(G)$

T: If G_n converges \Rightarrow — μ_{G_n} weakly converges
 — $\mu_{G_n}(x^2)$ converges

$$\int x^k d\mu_{G_n} = \frac{1}{|G_n|} \sum \lambda_i^k = \frac{1}{|G_n|} \text{tr}(A^k)$$

P : graph polynomial

P

- char. polys ✓

M_p : uniform measure on the roots of $p(G)$ in \mathbb{C}

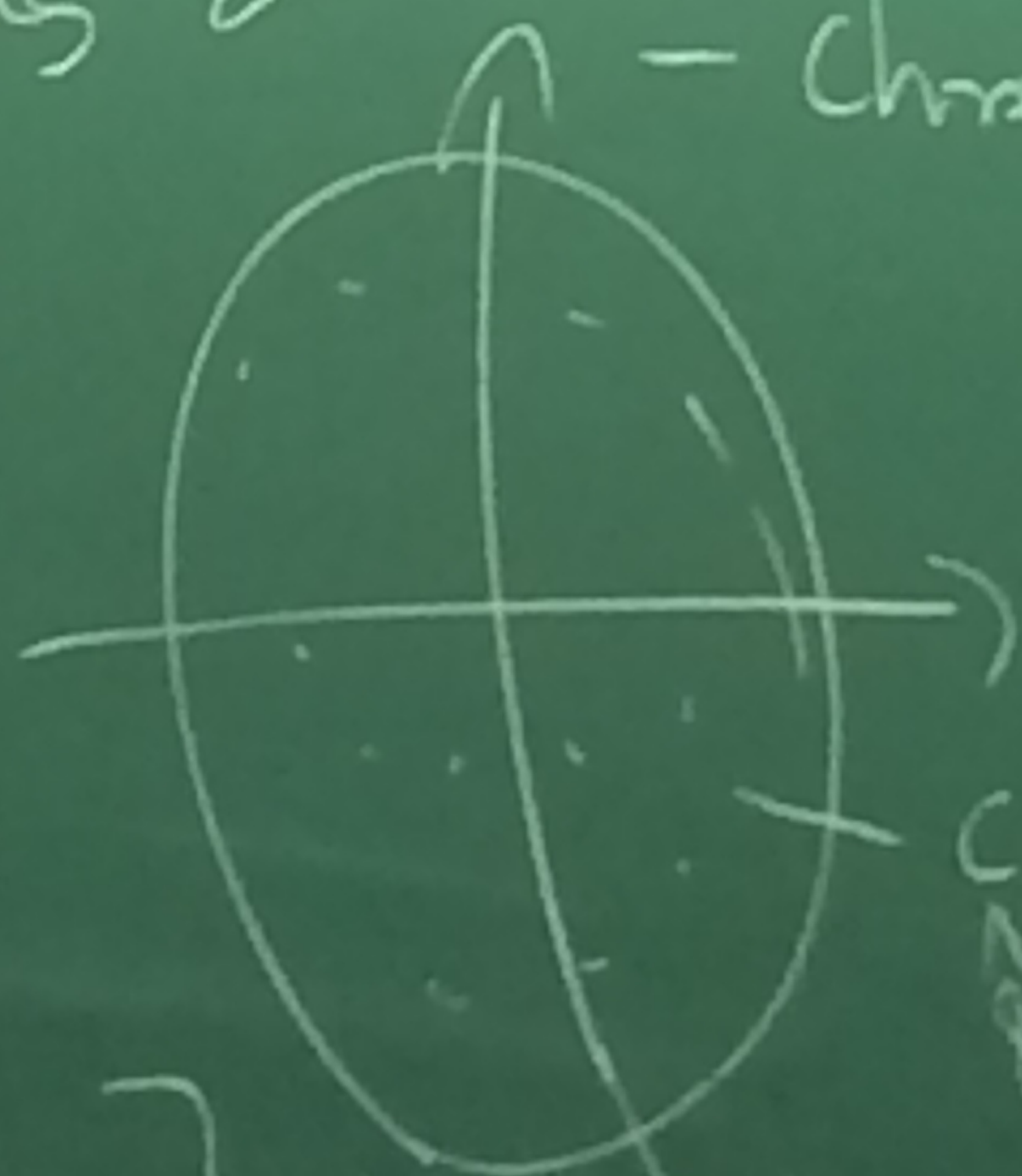
- Chromatic polynomial

of legal colorings of G

$\exists \chi_G$

$\chi_G(c) \equiv \chi_G$

χ_G



$\mathbb{C} \cap \mathbb{D}$ [Schubert]

S_G

T [A. Hubai]

If G_n BS-convex \Rightarrow

$S_{2^k} \text{ dgs. convex}$

[Zhang-Hubai] If G_n BS-converges $\Rightarrow \sum 2^k \text{dgs}_n$ converges

[Cribari-Frenkel] Generalized for a huge natural class of graph polynomials.

Matching polynomial.

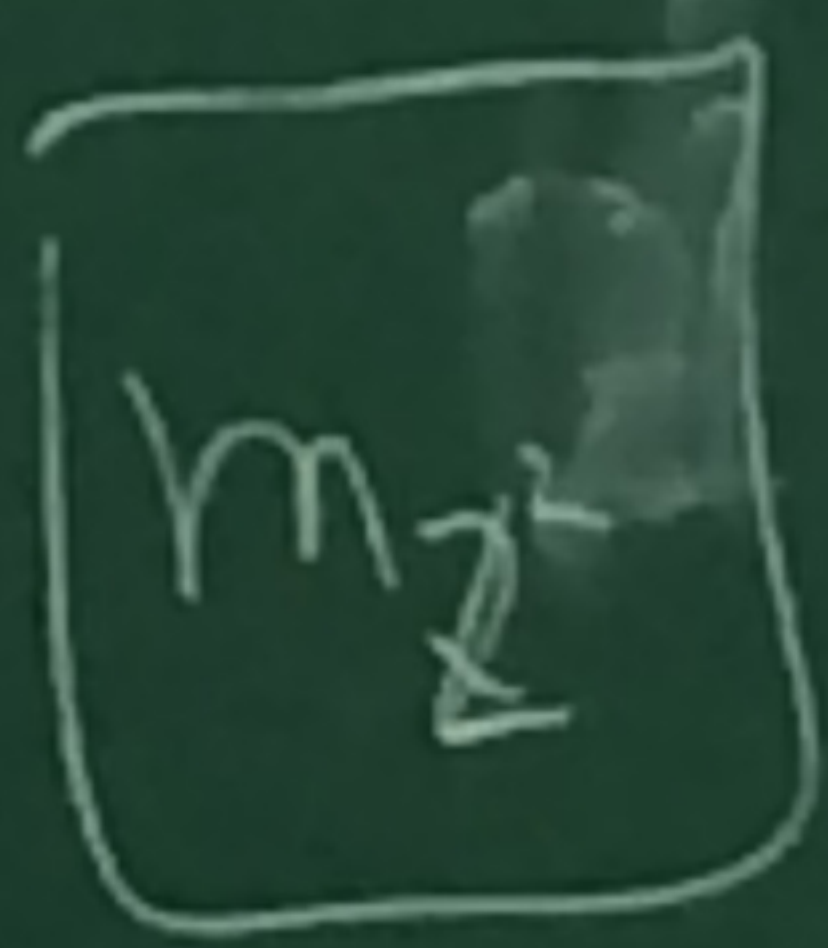
$$\sum (-1)^t m_t(G) x^t$$

$m_t(G)$ - matching measure
 $m_t(G)$ # of matchings of t edges in G



BS - arbitrary

HOWEVER, OVER ENTROPY!



WHAT IS IT?
 NO ATOMS

